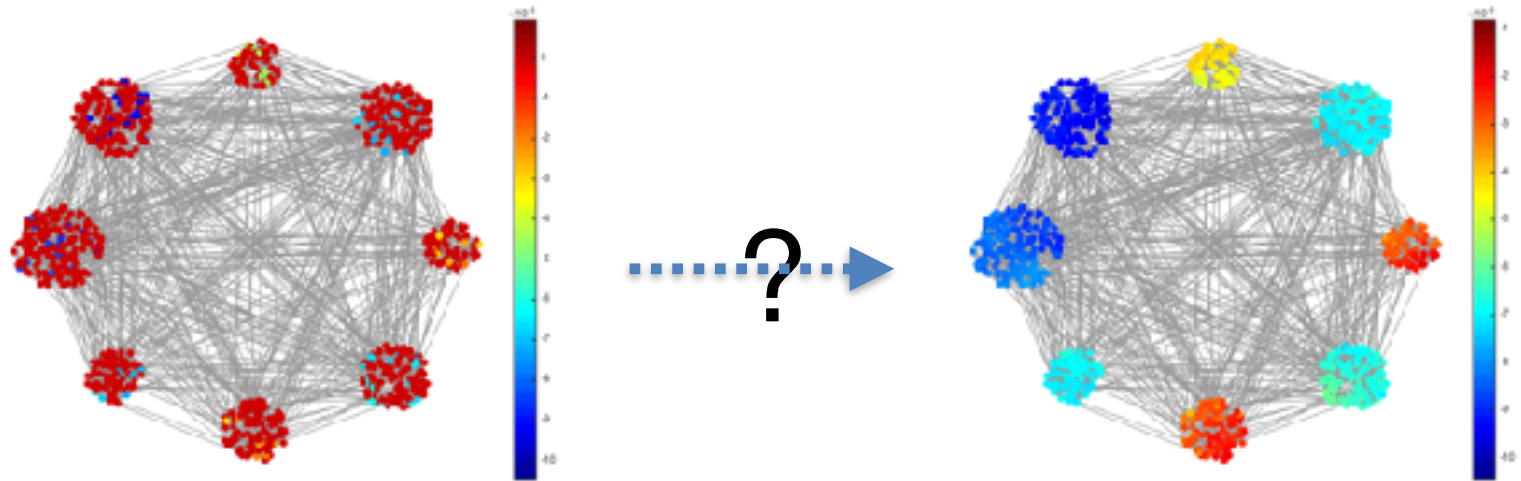


Sampling, Inference and Clustering for Data on Graphs

**Joint work with Gilles Puy, Nicolas Tremblay
and Rémi Gribonval**

Goal

Given partially observed information at the nodes of a graph



Can we robustly and efficiently infer missing information ?

What signal model ?

How many observations ?

Influence of the structure of the graph ?

Notations

$\mathcal{G} = \{\mathcal{V}, \mathcal{E}, W\}$ weighted, undirected

\mathcal{V} is the set of n nodes

\mathcal{E} is the set of edges

$W \in \mathbb{R}^{n \times n}$ is the weighted adjacency matrix

$L \in \mathbb{R}^{n \times n}$

combinatorial graph Laplacian $L := D - W$

normalised Laplacian $L := I - D^{-1/2} W D^{-1/2}$

diagonal degree matrix D has entries $d_i := \sum_{i \neq j} W_{ij}$

SP on Graphs Cheat Sheet

$\mathbf{x} \in \mathbb{R}^n$ a (scalar valued) signal

$$\mathbf{L} \in \mathbb{R}^{n \times n}$$

$$\mathbf{L} = \mathbf{U}\mathbf{\Lambda}\mathbf{U}^\top$$

Graph Fourier

Frequencies

Laplacian



$$-\Delta$$



$$-\Delta = \mathcal{F}^\top \omega^2 \mathcal{F}$$

Filter and Filtering

$$g(\mathbf{L}) = \mathbf{U}g(\mathbf{\Lambda})\mathbf{U}^\top$$

$$g(\mathbf{L})\mathbf{x}$$



$$g \star \mathbf{x}$$

$$\widehat{g(\mathbf{L})\mathbf{x}}(k) = g(\lambda_k)\hat{\mathbf{x}}(k)$$



$$\widehat{g \star \mathbf{x}}(\omega) = \hat{g}(\omega)\hat{\mathbf{x}}(\omega)$$

Notations

L is real, symmetric PSD

orthonormal eigenvectors $U \in \mathbb{R}^{n \times n}$ Graph Fourier Matrix

non-negative eigenvalues $\lambda_1 \leq \lambda_2 \leq \dots, \lambda_n$

$$L = U\Lambda U^T$$

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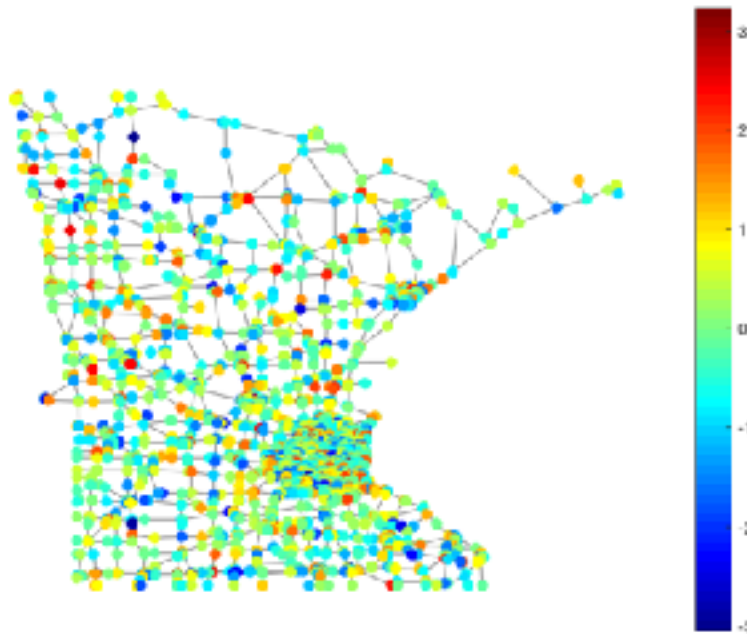
$$L = U\Lambda U^T$$

k -bandlimited signals $\mathbf{x} \in \mathbb{R}^n$

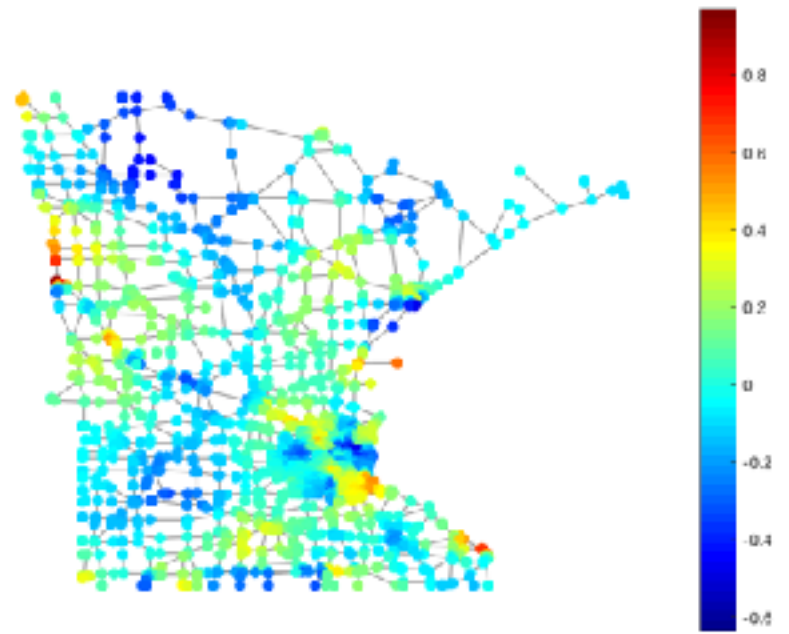
Fourier coefficients $\hat{\mathbf{x}} = U^T \mathbf{x}$

$$\mathbf{x} = U_k \hat{\mathbf{x}}^k \quad \hat{\mathbf{x}}^k \in \mathbb{R}^k$$

$U_k := (\mathbf{u}_1, \dots, \mathbf{u}_k) \in \mathbb{R}^{n \times k}$ first k eigenvectors only



white noise



band-limited approximation

Sampling Model

$$\mathbf{p} \in \mathbb{R}^n \quad \mathbf{p}_i > 0 \quad \|\mathbf{p}\|_1 = \sum_{i=1}^n \mathbf{p}_i = 1$$

$$\mathbf{P} := \text{diag}(\mathbf{p}) \in \mathbb{R}^{n \times n}$$

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Draw independently m samples (random sampling)

$$\mathbb{P}(\omega_j = i) = \mathbf{p}_i, \quad \forall j \in \{1, \dots, m\} \text{ and } \forall i \in \{1, \dots, n\}$$

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$$\mathbf{y}_j := \mathbf{x}_{\omega_j}, \quad \forall j \in \{1, \dots, m\}$$

$$\mathbf{y} = \mathbf{M}\mathbf{x}$$

Sampling Model

$$\frac{\|\mathbf{U}_k^\top \boldsymbol{\delta}_i\|_2}{\|\mathbf{U}^\top \boldsymbol{\delta}_i\|_2} = \frac{\|\mathbf{U}_k^\top \boldsymbol{\delta}_i\|_2}{\|\boldsymbol{\delta}_i\|_2} = \|\mathbf{U}_k^\top \boldsymbol{\delta}_i\|_2$$

How much a perfect impulse can be concentrated on first k eigenvectors

Carries interesting information about the graph

Sampling Model

$$\frac{\|\mathbf{U}_k^T \boldsymbol{\delta}_i\|_2}{\|\mathbf{U}^T \boldsymbol{\delta}_i\|_2} = \frac{\|\mathbf{U}_k^T \boldsymbol{\delta}_i\|_2}{\|\boldsymbol{\delta}_i\|_2} = \|\mathbf{U}_k^T \boldsymbol{\delta}_i\|_2$$

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Ideally: p_i large wherever $\|\mathbf{U}_k^T \boldsymbol{\delta}_i\|_2$ is large

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Graph Coherence

$$\nu_{\mathbf{p}}^k := \max_{1 \leq i \leq n} \left\{ \mathbf{p}_i^{-1/2} \|\mathbf{U}_k^T \boldsymbol{\delta}_i\|_2 \right\}$$

$$\text{Rem: } \nu_{\mathbf{p}}^k \geq \sqrt{k}$$

Stable Embedding

Theorem 1 (Restricted isometry property). *Let \mathbf{M} be a random subsampling matrix with the sampling distribution \mathbf{p} . For any $\delta, \epsilon \in (0, 1)$, with probability at least $1 - \epsilon$,*

$$(1 - \delta) \|\mathbf{x}\|_2^2 \leq \frac{1}{m} \left\| \mathbf{M} \mathbf{P}^{-1/2} \mathbf{x} \right\|_2^2 \leq (1 + \delta) \|\mathbf{x}\|_2^2 \quad (1)$$

for all $\mathbf{x} \in \text{span}(\mathbf{U}_k)$ provided that

$$m \geq \frac{3}{\delta^2} (\nu_{\mathbf{p}}^k)^2 \log \left(\frac{2k}{\epsilon} \right). \quad (2)$$

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$\mathbf{M} \mathbf{P}^{-1/2} \mathbf{x} = \mathbf{P}_{\Omega}^{-1/2} \mathbf{M} \mathbf{x}$ Only need \mathbf{M} , re-weighting offline

Stable Embedding

Theorem 1 (Restricted isometry property). *Let M be a random subsampling matrix with the sampling distribution \mathbf{p} . For any $\delta, \epsilon \in (0, 1)$, with probability at least $1 - \epsilon$,*

$$(1 - \delta) \|\mathbf{x}\|_2^2 \leq \frac{1}{m} \left\| MP^{-1/2} \mathbf{x} \right\|_2^2 \leq (1 + \delta) \|\mathbf{x}\|_2^2 \quad (1)$$

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$$MP^{-1/2} \mathbf{x} = P_{\Omega}^{-1/2} M\mathbf{x}$$

Only need M , re-weighting offline

$$(\nu_{\mathbf{p}}^k)^2 \geq k$$

Need to sample at least k nodes

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$(\nu_{\mathbf{p}}^k)^2 \geq k$ Need to sample at least k nodes

Proof similar to CS in bounded ONB but simpler since model is a subspace (not a union)

Stable Embedding

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Can we reduce to optimal amount ?

Variable Density Sampling $\mathbf{p}_i^* := \frac{\|\mathbf{U}_k^T \boldsymbol{\delta}_i\|_2^2}{k}, \quad i = 1, \dots, n$

is such that: $(\nu_{\mathbf{p}}^k)^2 = k$ and depends on structure of graph

Stable Embedding

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Corollary 1. *Let \mathbf{M} be a random subsampling matrix constructed with the sampling distribution \mathbf{p}^* . For any $\delta, \epsilon \in (0, 1)$, with probability at least $1 - \epsilon$,*

$$(1 - \delta) \|\mathbf{x}\|_2^2 \leq \frac{1}{m} \left\| \mathbf{M} \mathbf{P}^{-1/2} \mathbf{x} \right\|_2^2 \leq (1 + \delta) \|\mathbf{x}\|_2^2$$

for all $\mathbf{x} \in \text{span}(\mathbf{U}_k)$ provided that

$$m \geq \frac{3}{\delta^2} k \log \left(\frac{2k}{\epsilon} \right).$$

Recovery Procedures (Inference)

$$\mathbf{y} = \mathbf{M}\mathbf{x} + \mathbf{n} \quad \mathbf{y} \in \mathbb{R}^m$$

$$\mathbf{x} \in \text{span}(\mathbf{U}_k) \quad \text{stable embedding}$$

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Standard Decoder

$$\min_{\mathbf{z} \in \text{span}(\mathbf{U}_k)} \left\| \mathbf{P}_{\Omega}^{-1/2} (\mathbf{M}\mathbf{z} - \mathbf{y}) \right\|_2$$

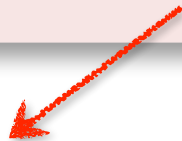
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need projector

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need projector

re-weighting for RIP

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Efficient Decoder:

$$\min_{\mathbf{z} \in \mathbb{R}^n} \left\| \mathbf{P}_{\Omega}^{-1/2} (\mathbf{M}\mathbf{z} - \mathbf{y}) \right\|_2^2 + \gamma \mathbf{z}^T g(\mathbf{L})\mathbf{z}$$

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soft constrain on frequencies
efficient implementation

Analysis of Standard Decoder

Standard Decoder:

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Theorem 1. *Let Ω be a set of m indices selected independently from $\{1, \dots, n\}$ with sampling distribution $\mathbf{p} \in \mathbb{R}^n$, and \mathbf{M} the associated sampling matrix. Let $\epsilon, \delta \in (0, 1)$ and $m \geq \frac{3}{\delta^2} (\nu_{\mathbf{p}}^k)^2 \log\left(\frac{2k}{\epsilon}\right)$. With probability at least $1 - \epsilon$, the following holds for all $\mathbf{x} \in \text{span}(\mathbf{U}_k)$ and all $\mathbf{n} \in \mathbb{R}^m$.*

i) *Let \mathbf{x}^* be the solution of Standard Decoder with $\mathbf{y} = \mathbf{M}\mathbf{x} + \mathbf{n}$. Then,*

$$\|\mathbf{x}^* - \mathbf{x}\|_2 \leq \frac{2}{\sqrt{m(1-\delta)}} \left\| \mathbf{P}_\Omega^{-1/2} \mathbf{n} \right\|_2. \quad (1)$$

ii) *There exist particular vectors $\mathbf{n}_0 \in \mathbb{R}^m$ such that the solution \mathbf{x}^* of Standard Decoder with $\mathbf{y} = \mathbf{M}\mathbf{x} + \mathbf{n}_0$ satisfies*

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Exact recovery when noiseless

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Analysis of Efficient Decoder

Efficient Decoder:

$$\min_{\mathbf{z} \in \mathbb{R}^n} \left\| \mathbf{P}_{\Omega}^{-1/2} (\mathbf{M}\mathbf{z} - \mathbf{y}) \right\|_2^2 + \gamma \mathbf{z}^T g(\mathbf{L})\mathbf{z}$$

non-negative

Analysis of Efficient Decoder

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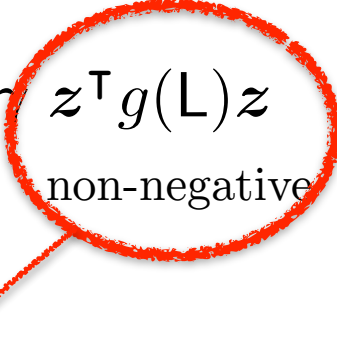
Filter reshapes Fourier coefficients

$$h : \mathbb{R} \rightarrow \mathbb{R} \quad \mathbf{x}_h := \mathbf{U} \text{diag}(\hat{\mathbf{h}}) \mathbf{U}^T \mathbf{x} \in \mathbb{R}^n$$

$$\hat{\mathbf{h}} = (h(\boldsymbol{\lambda}_1), \dots, h(\boldsymbol{\lambda}_n))^T \in \mathbb{R}^n$$

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$$p(t) = \sum_{i=0}^d \alpha_i t^i \quad \mathbf{x}_p = \mathbf{U} \text{diag}(\hat{\mathbf{p}}) \mathbf{U}^T \mathbf{x} = \sum_{i=0}^d \alpha_i \mathbf{L}^i \mathbf{x}$$

Pick special polynomials and use e.g. recurrence relations for fast filtering (with sparse matrix-vector multiply only)

Analysis of Efficient Decoder

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non-decreasing = 
penalizes high-frequencies

Analysis of Efficient Decoder

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non-decreasing = 
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Favours reconstruction of approximately band-limited signals

Ideal filter yields Standard Decoder

$$i_{\lambda_k}(t) := \begin{cases} 0 & \text{if } t \in [0, \lambda_k], \\ +\infty & \text{otherwise,} \end{cases}$$

Analysis of Efficient Decoder

Theorem 1. *Let Ω , M , P , m as before and $M_{\max} > 0$ be a constant such that $\|MP^{-1/2}\|_2 \leq M_{\max}$. Let $\epsilon, \delta \in (0, 1)$. With probability at least $1 - \epsilon$, the following holds for all $\mathbf{x} \in \text{span}(\mathbf{U}_k)$, all $\mathbf{n} \in \mathbb{R}^n$, all $\gamma > 0$, and all nonnegative and nondecreasing polynomial functions g such that $g(\boldsymbol{\lambda}_{k+1}) > 0$.*

Let \mathbf{x}^ be the solution of Efficient Decoder with $\mathbf{y} = M\mathbf{x} + \mathbf{n}$. Then,*

$$\|\boldsymbol{\alpha}^* - \mathbf{x}\|_2 \leq \frac{1}{\sqrt{m(1-\delta)}} \left[\left(2 + \frac{M_{\max}}{\sqrt{\gamma g(\boldsymbol{\lambda}_{k+1})}} \right) \left\| P_{\Omega}^{-1/2} \mathbf{n} \right\|_2 + \left(M_{\max} \sqrt{\frac{g(\boldsymbol{\lambda}_k)}{g(\boldsymbol{\lambda}_{k+1})}} + \sqrt{\gamma g(\boldsymbol{\lambda}_k)} \right) \|\mathbf{x}\|_2 \right], \quad (1)$$

and

$$\|\boldsymbol{\beta}^*\|_2 \leq \frac{1}{\sqrt{\gamma g(\boldsymbol{\lambda}_{k+1})}} \left\| P_{\Omega}^{-1/2} \mathbf{n} \right\|_2 + \sqrt{\frac{g(\boldsymbol{\lambda}_k)}{g(\boldsymbol{\lambda}_{k+1})}} \|\mathbf{x}\|_2, \quad (2)$$

where $\boldsymbol{\alpha}^* := \mathbf{U}_k \mathbf{U}_k^T \mathbf{x}^*$ and $\boldsymbol{\beta}^* := (\mathbf{I} - \mathbf{U}_k \mathbf{U}_k^T) \mathbf{x}^*$.

Analysis of Efficient Decoder

Noiseless case:

$$\|\mathbf{x}^* - \mathbf{x}\|_2 \leq \frac{1}{\sqrt{m(1-\delta)}} \left(M_{\max} \sqrt{\frac{g(\boldsymbol{\lambda}_k)}{g(\boldsymbol{\lambda}_{k+1})}} + \sqrt{\gamma g(\boldsymbol{\lambda}_k)} \right) \|\mathbf{x}\|_2 + \sqrt{\frac{g(\boldsymbol{\lambda}_k)}{g(\boldsymbol{\lambda}_{k+1})}} \|\mathbf{x}\|_2$$

$g(\boldsymbol{\lambda}_k) = 0$ + non-decreasing implies perfect reconstruction

Analysis of Efficient Decoder

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$g(\boldsymbol{\lambda}_k) = 0$ + non-decreasing implies perfect reconstruction

Otherwise:

choose γ as close as possible to 0 and seek to minimise the ratio $g(\boldsymbol{\lambda}_k)/g(\boldsymbol{\lambda}_{k+1})$

Choose filter to increase spectral gap ?

Clusters are of course good

Noise: $\|\mathbf{P}_\Omega^{-1/2} \mathbf{n}\|_2 / \|\mathbf{x}\|_2$

Estimating the Optimal Distribution

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Need to estimate $\|U_k^T \delta_i\|_2^2$

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Need to estimate $\|\mathbf{U}_k^T \boldsymbol{\delta}_i\|_2^2$

Filter random signals with ideal low-pass filter:

$$\mathbf{r}_{b_{\lambda_k}} = \mathbf{U} \text{diag}(\boldsymbol{\lambda}_1, \dots, \boldsymbol{\lambda}_k, 0, \dots, 0) \mathbf{U}^T \mathbf{r} = \mathbf{U}_k \mathbf{U}_k^T \mathbf{r}$$

$$\mathbb{E}(\mathbf{r}_{b_{\lambda_k}})_i^2 = \boldsymbol{\delta}_i^T \mathbf{U}_k \mathbf{U}_k^T \mathbb{E}(\mathbf{r} \mathbf{r}^T) \mathbf{U}_k \mathbf{U}_k^T \boldsymbol{\delta}_i = \|\mathbf{U}_k^T \boldsymbol{\delta}_i\|_2^2$$

Estimating the Optimal Distribution

Need to estimate $\|\mathbf{U}_k^\top \boldsymbol{\delta}_i\|_2^2$

Filter random signals with ideal low-pass filter:

$$\mathbf{r}_{b_{\lambda_k}} = \mathbf{U} \operatorname{diag}(\boldsymbol{\lambda}_1, \dots, \boldsymbol{\lambda}_k, 0, \dots, 0) \mathbf{U}^\top \mathbf{r} = \mathbf{U}_k \mathbf{U}_k^\top \mathbf{r}$$

$$\mathbb{E}(\mathbf{r}_{b_{\lambda_k}})_i^2 = \boldsymbol{\delta}_i^\top \mathbf{U}_k \mathbf{U}_k^\top \mathbb{E}(\mathbf{r} \mathbf{r}^\top) \mathbf{U}_k \mathbf{U}_k^\top \boldsymbol{\delta}_i = \|\mathbf{U}_k^\top \boldsymbol{\delta}_i\|_2^2$$

In practice, one may use a polynomial approximation of the ideal filter and:

$$\tilde{\mathbf{p}}_i := \frac{\sum_{l=1}^L (\mathbf{r}_{c_{\lambda_k}}^l)_i^2}{\sum_{i=1}^n \sum_{l=1}^L (\mathbf{r}_{c_{\lambda_k}}^l)_i^2}$$

$$L \geq \frac{C}{\delta^2} \log \left(\frac{2n}{\epsilon} \right)$$

Estimating the Eigengap

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Again, low-pass filtering random signals:

$$(1 - \delta) \sum_{i=1}^n \left\| \mathbf{U}_{j^*}^\top \boldsymbol{\delta}_i \right\|_2^2 \leq \sum_{i=1}^n \sum_{l=1}^L (\mathbf{r}_{b_\lambda}^l)_i^2 \leq (1 + \delta) \sum_{i=1}^n \left\| \mathbf{U}_{j^*}^\top \boldsymbol{\delta}_i \right\|_2^2$$

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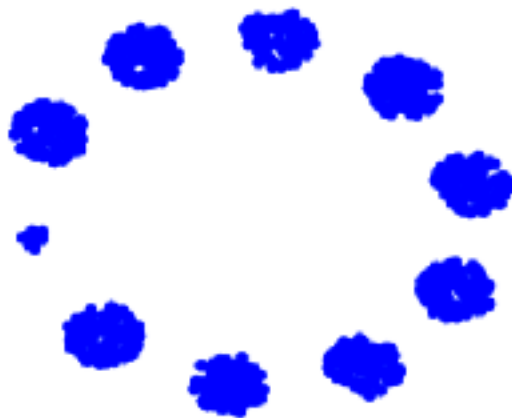
Since:
$$\sum_{i=1}^n \|\mathbf{U}_{j^*}^\top \boldsymbol{\delta}_i\|_2^2 = \|\mathbf{U}_{j^*}\|_{\text{Frob}}^2 = j^*$$

We have:
$$(1 - \delta) j^* \leq \sum_{i=1}^n \sum_{l=1}^L (\mathbf{r}_{b_\lambda}^l)_i^2 \leq (1 + \delta) j^*$$

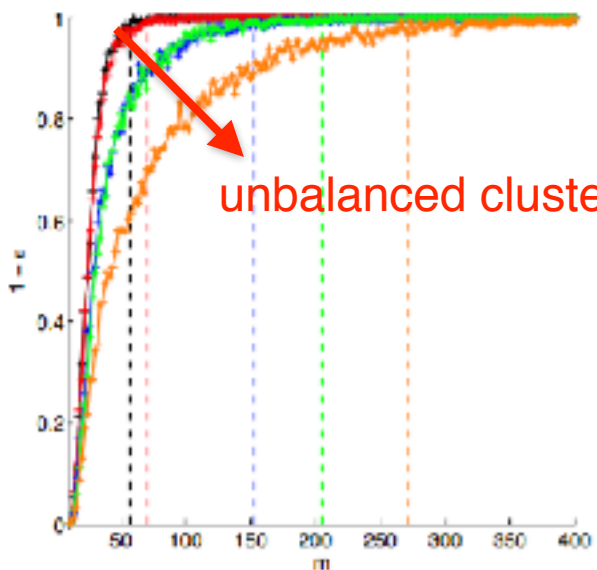
Dichotomy using the filter bandwidth

Experiments

Community graph

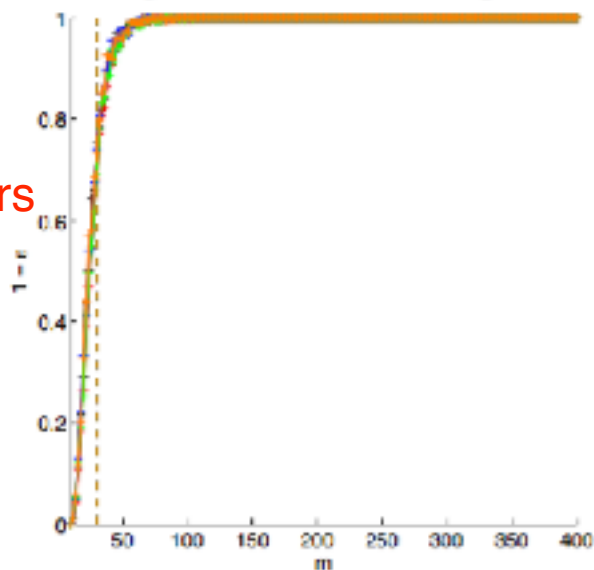


Uniform distribution π

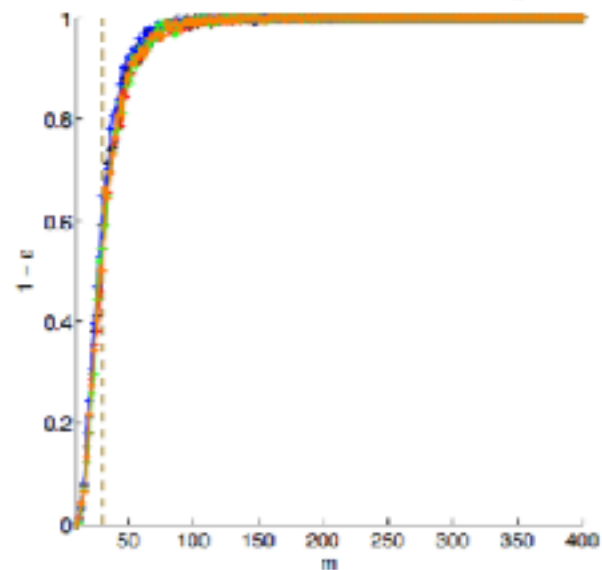


unbalanced clusters

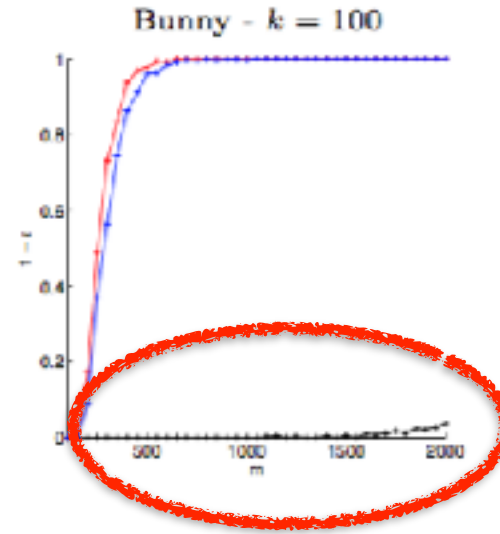
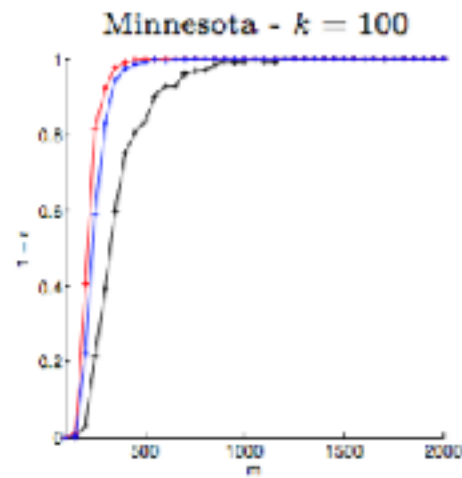
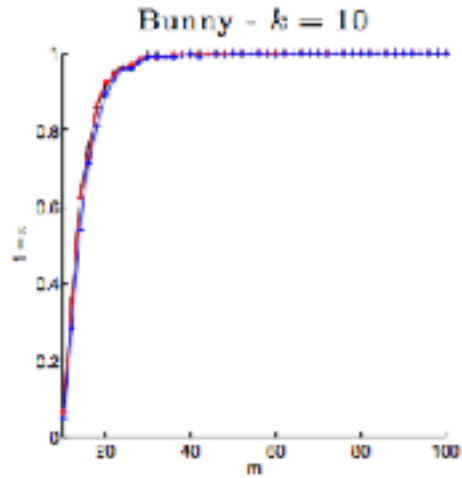
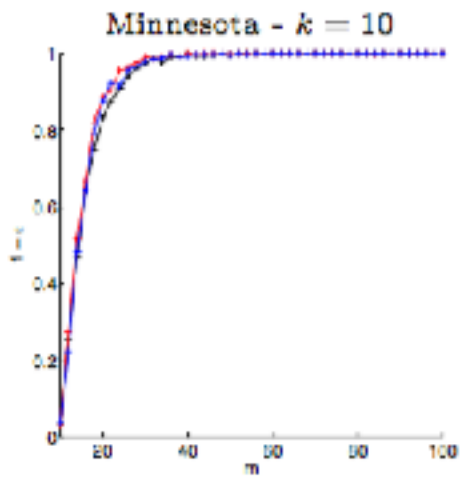
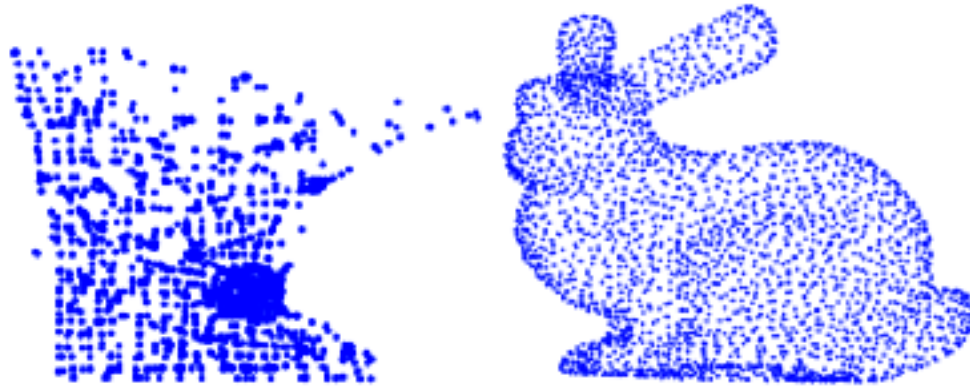
Optimal distribution p^*



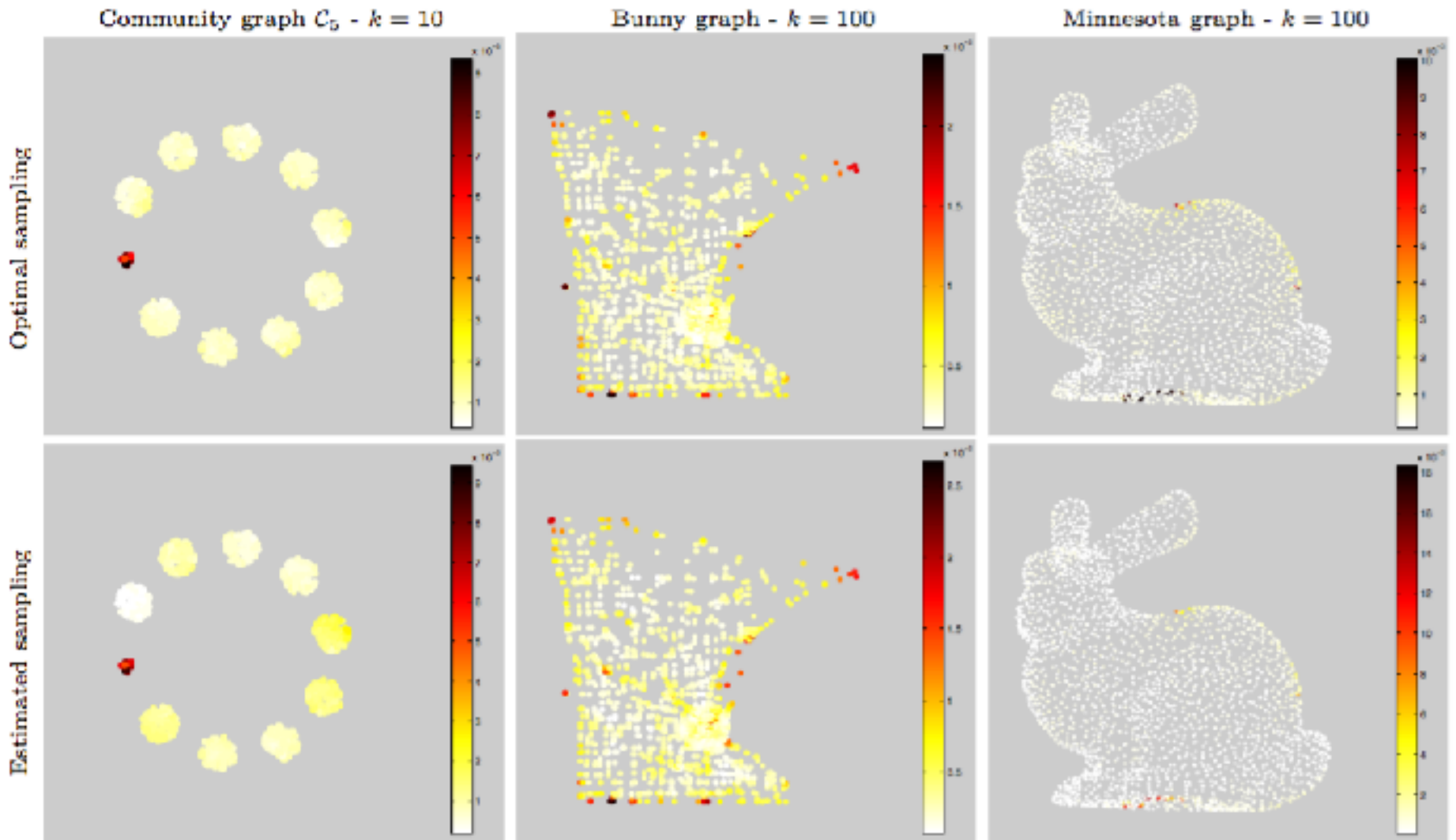
Estimated distribution \bar{p}



Experiments



Experiments

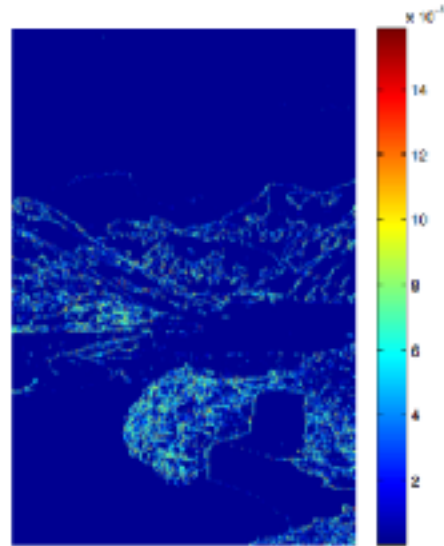


Experiments



(a)

Original



(b)

Reconstructed (sampling with $\tilde{\mathbf{y}}$)



7%

Compressive Spectral Clustering

Clustering equivalent to recovery of cluster assignment functions

Well-defined clusters \rightarrow band-limited assignment functions!

Compressive Spectral Clustering

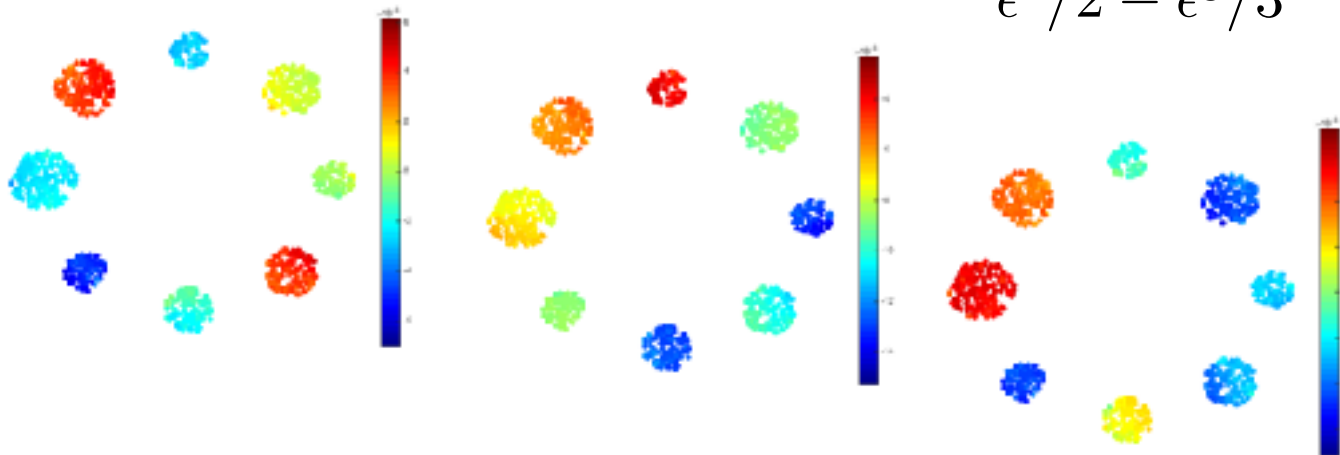
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Generate features by filtering random signals

by Johnson-Lindenstrauss

$$\eta = \frac{4 + 2\beta}{\epsilon^2/2 - \epsilon^3/3} \log n$$



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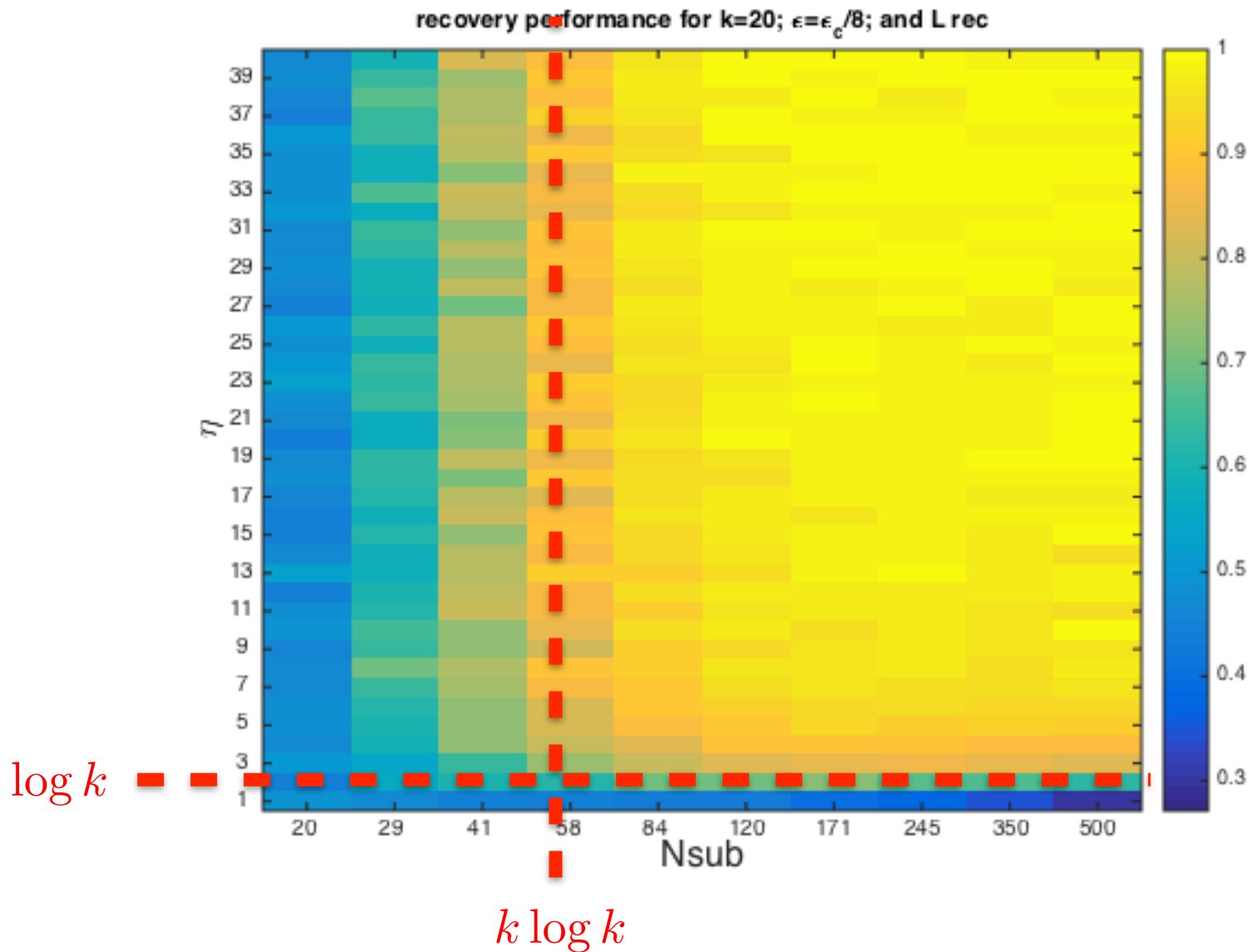
by Johnson-Lindenstrauss
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Each feature map is smooth, therefore keep

$$m \geq \frac{6}{\delta^2} \nu_k^2 \log \left(\frac{k}{\epsilon'} \right)$$

Use k-means on compressed data and feed into Efficient Decoder²⁵

Compressive Spectral Clustering



Conclusion

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- ***Stable, robust and universal random sampling*** of smoothly varying information on graphs.

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- ***Stable, robust and universal random sampling*** of smoothly varying information on graphs.
- Tractable decoder with guarantees
- ***Optimal sampling distribution*** depends on graph structure
- Can be used for inference, (SVD less) compressive clustering

Thank you !