

Sparse recovery with (non) convex optimization

Rémi Gribonval

METISS project-team (audio signal processing, speech recognition, source separation)
INRIA, Rennes, France



IRISA

UNE UNITÉ DE RECHERCHE À LA POINTE DES SCIENCES
ET DES TECHNOLOGIES DE L'INFORMATION
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Sparsity and Inverse Problems

- More unknowns than equations:

$$\mathbf{A}x_0 = \mathbf{A}x_1 \not\Rightarrow x_0 = x_1$$

- Uniqueness of sparse solutions = identifiability

- ♦ if x_0, x_1 are “sufficiently sparse”,
- ♦ then $\mathbf{A}x_0 = \mathbf{A}x_1 \Rightarrow x_0 = x_1$

- If x_0, x_1 “sufficiently sparse”, identification with

- ♦ L1-minimization = convex problem
- ♦ Greedy algorithms

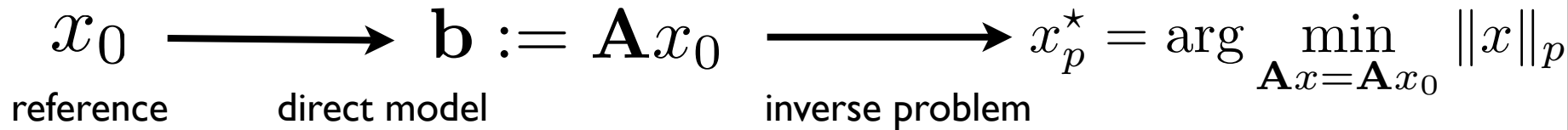
- Robustness to “approximately sparse” and noise

Overview

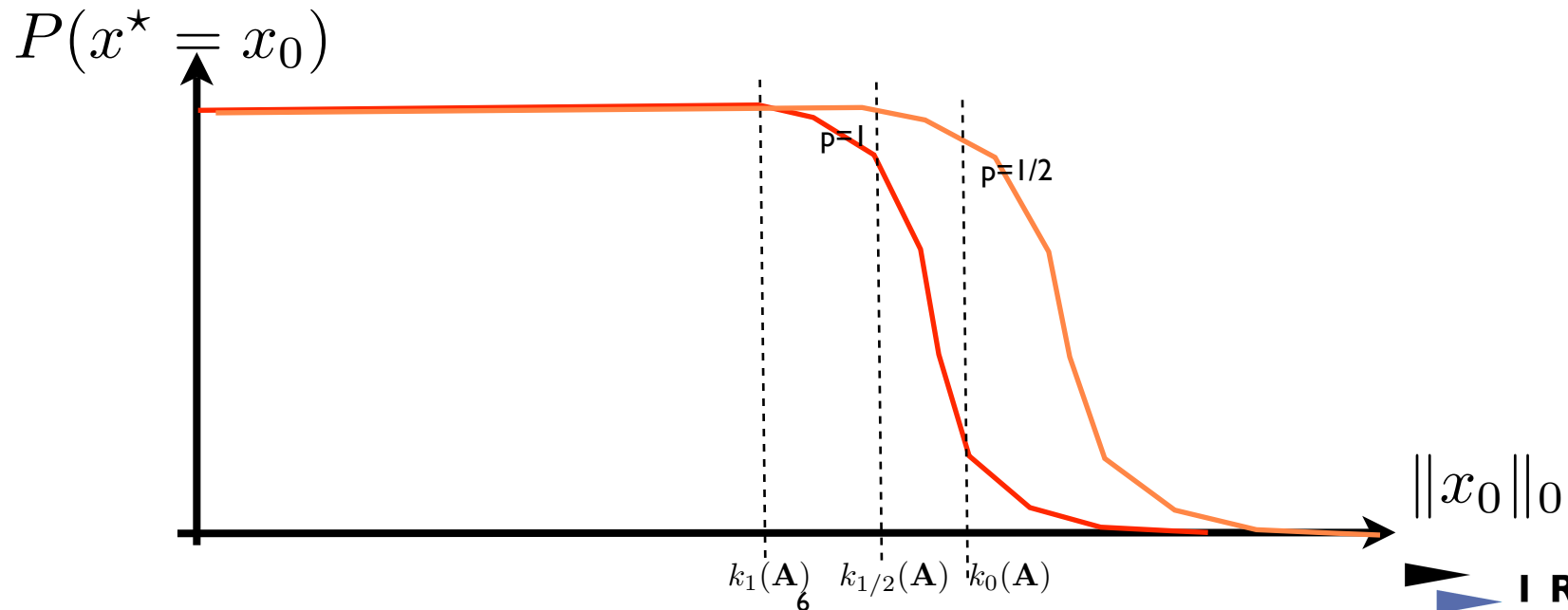
- Exact recovery
 - ◆ greedy algorithms
 - ◆ (non)convex L_p -minimization
 - ◆ comparisons
- Stability and robustness
 - ◆ Instance optimality for L_p -minimization
 - ◆ Restricted Isometry Property

Exact recovery

Empirical observation : L_p versus L_1



Typical observation (e.g. Chartrand 2007) + extrapolation



Proved Equivalence between L0 and L1

- “Empty” theorem : assume that $\mathbf{b} = \mathbf{A}x_0$
 - ◆ if $\|x_0\|_0 \leq k_0(\mathbf{A})$ then $x_0 = x_0^*$
 - ◆ if $\|x_0\|_0 \leq k_1(\mathbf{A})$ then $x_0 = x_1^*$

- Content = estimation of $k_0(\mathbf{A})$ and $k_1(\mathbf{A})$
 - ◆ Donoho & Huo 2001 : *pair of bases, coherence*
 - ◆ Donoho & Elad 2003, Gribonval & Nielsen 2003 : *dictionary, coherence*
 - ◆ Candes, Romberg, Tao 2004 : *random dictionaries,* *restricted isometry constants*
 - ◆ Tropp 2004 : *idem for Orthonormal Matching Pursuit,* *cumulative coherence*

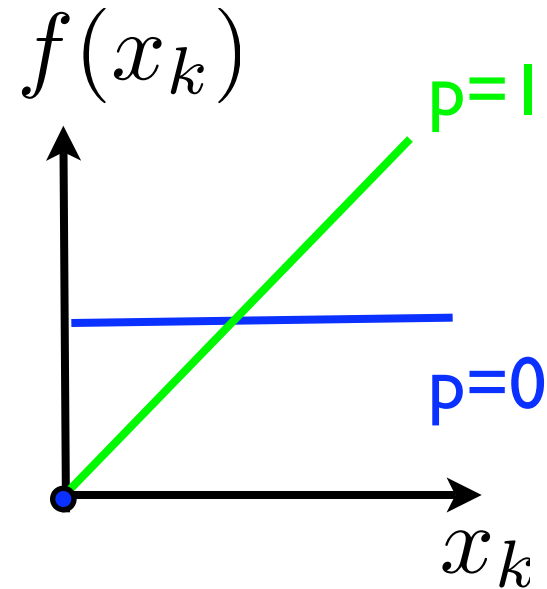
- What about $x_p^*, 0 \leq p \leq 1$?

General sparsity measures

- Lp-norms $\|x\|_p^p := \sum_k |x_k|^p, 0 \leq p \leq 1$

- f-norms! $\|x\|_f := \sum_k f(|x_k|)$

- Constrained minimization

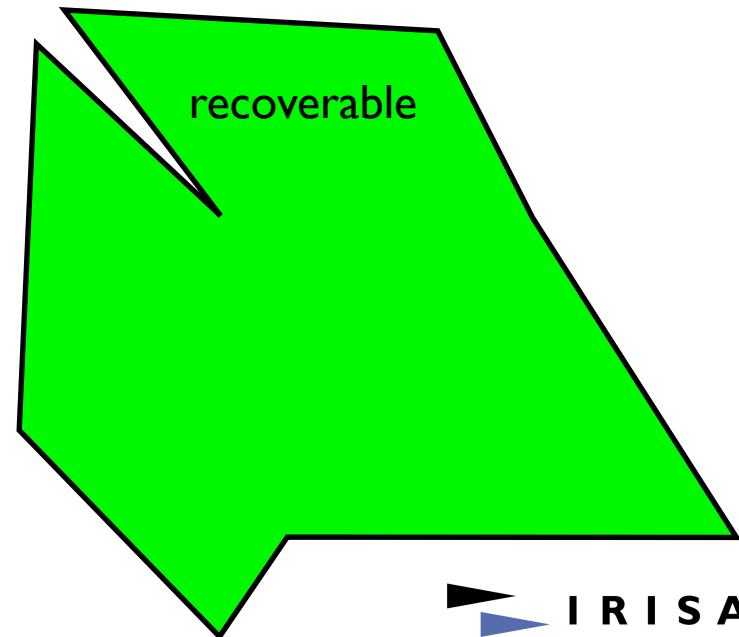


$$x_f^* = x_f^*(\mathbf{b}, \mathbf{A}) \in \arg \min_x \|x\|_f \quad \text{subject to} \quad \mathbf{b} = \mathbf{A}x$$

When do we have $x_f^*(\mathbf{A}x_0, \mathbf{A}) = x_0$?

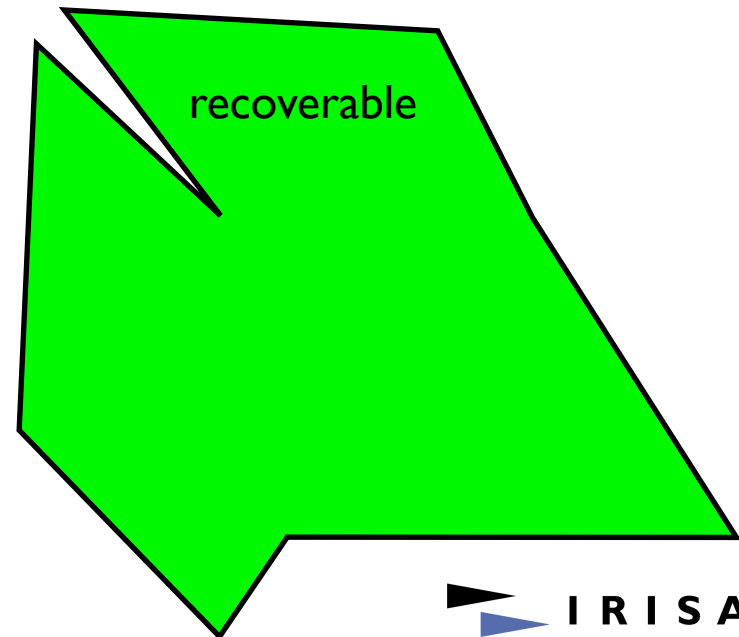
Recovery analysis for inverse problem $\mathbf{b} = \mathbf{A}x$

- Holy grail = characterize set of “recoverable” coefficients
$$\{x \in \mathbb{R}^N, x_f^*(\mathbf{A}x, \mathbf{A}) = x\}$$



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- Practice : find “simple” recovery conditions



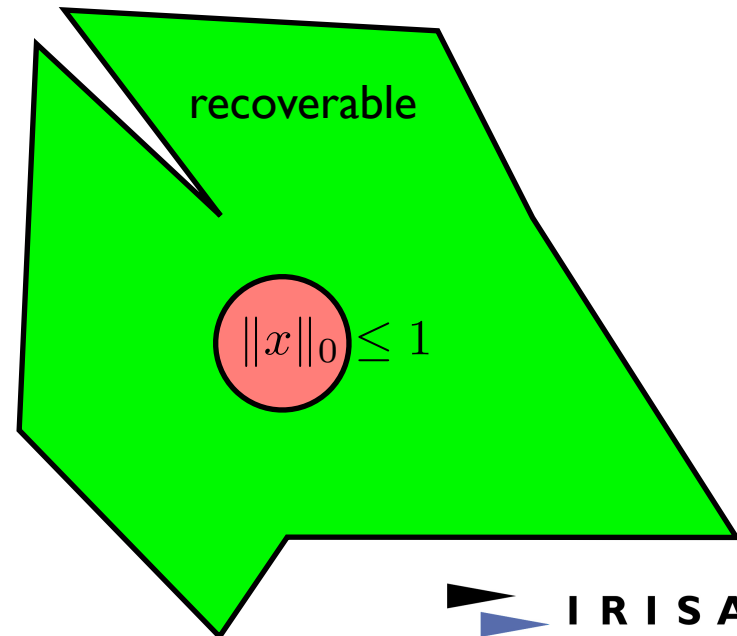
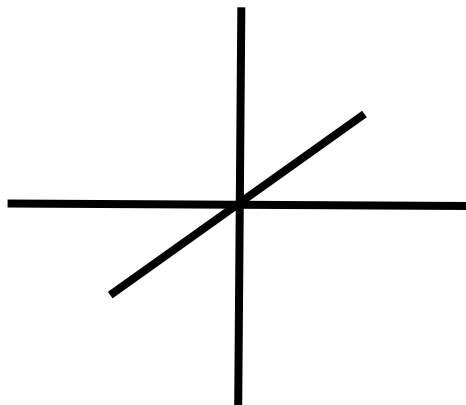
Recovery analysis for inverse problem $b = Ax$

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$$\{x \in \mathbb{R}^N, x_f^*(Ax, A) = x\}$$

- Practice : find “simple” recovery conditions

- ♦ l_1 -sparse



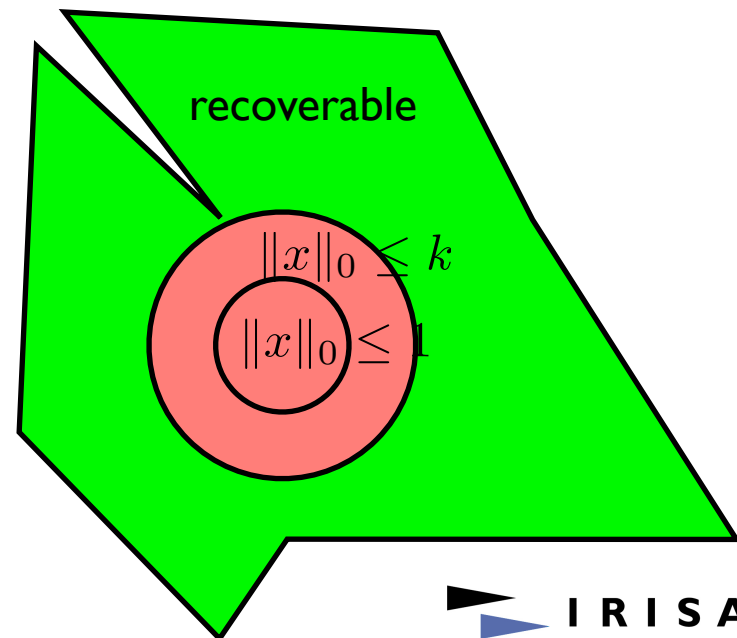
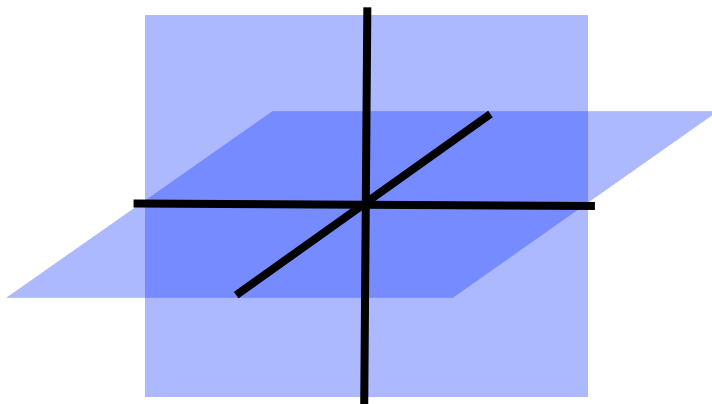
Recovery analysis for inverse problem $b = Ax$

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- Practice : find “simple” recovery conditions

- ◆ 1-sparse
- ◆ 2-sparse



Some “simple” recovery conditions

Support

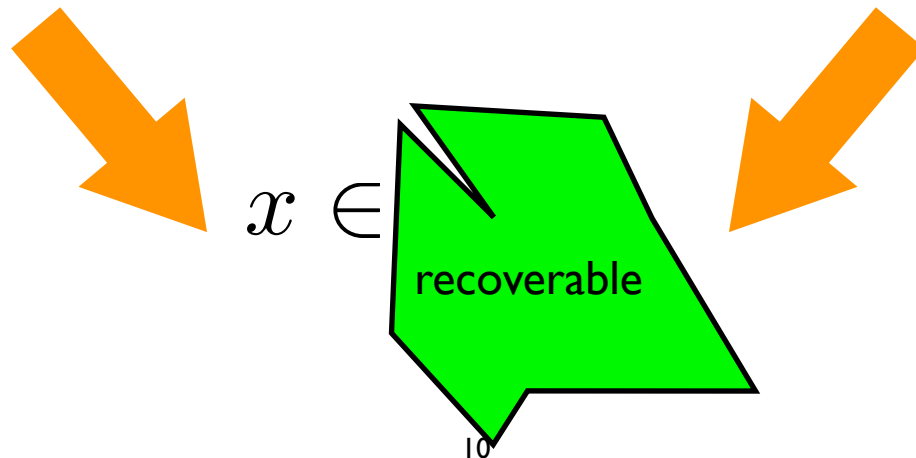
Sparsity level

“recoverable supports” =
subsets $I \subset \llbracket 1, N \rrbracket$
such that

$$\text{supp}(x) := \{k, x_k \neq 0\} \subset I$$

“recoverable sparsity” =
integers k
such that

$$\|x_0\|_0 \leq k$$



Exact recovery: Greedy algorithms

Objective

- Given
 - ◆ Dictionary \mathbf{A} ,
 - ◆ support set I
- Goal = understand when MP (and variants) is guaranteed to only select atoms in I , given any input vector $\mathbf{b} = \mathbf{A}_I \mathbf{x}_I$
- Analysis: based on *operator norms*

Operator norms

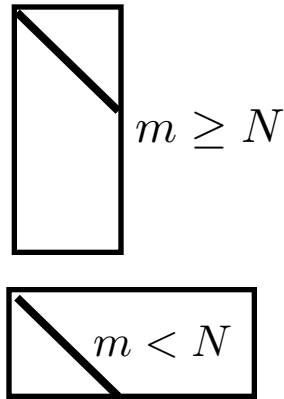
- Linear operator $\mathbf{L} : \ell^p \rightarrow \ell^q$
- Operator norm $\|\mathbf{L}\|_{p \rightarrow q} = \sup_{x \neq 0} \frac{\|\mathbf{L}x\|_q}{\|x\|_p}$
- Adjoint operator $\forall x, y \langle \mathbf{L}x, y \rangle = \langle x, \mathbf{L}^*y \rangle$
- For real-valued matrices $\mathbf{L}^* = \mathbf{L}^T$
- Duality relation: for all \mathbf{L} , $1 \leq p, q \leq \infty$

$$\|\mathbf{L}\|_{p \rightarrow q} = \|\mathbf{L}^T\|_{q' \rightarrow p'} \quad \frac{1}{p} + \frac{1}{p'} = 1; \quad \frac{1}{q} + \frac{1}{q'} = 1$$

Operator norms (ctd)

- When $p=q=2$
 - ♦ Singular Value Decomposition (SVD)

$$\mathbf{L} = U\Sigma V$$
$$UU^T = U^T U = Id_m$$
$$VV^T = V^T V = Id_N$$
$$\Sigma = \text{diag}(\sigma_i), \sigma_1 \geq \dots \geq 0$$



- ♦ Operator norm

$$\|\mathbf{L}\|_{2 \rightarrow 2} = \|\Sigma\|_{2 \rightarrow 2} = \sigma_{\max}(\mathbf{L})$$

- ♦ Proof: exercise

Operator norms (ctd)

- When $p=1$, for any q

- ♦ Columns

$$\mathbf{L} = [\mathbf{L}_1, \dots, \mathbf{L}_N]$$

- ♦ Operator norm

$$\|\mathbf{L}\|_{1 \rightarrow q} = \max_n \|\mathbf{L}_n\|_q$$

- ♦ Proof: exercise

Exact Recovery Condition for *MP

- **Theorem:** consider any weak/stagewise greedy algorithm that iterates

- ✦ a **selection** of atoms in a set Γ_i such that

$$\inf_{l \in \Gamma_i} |\mathbf{A}_l^T \mathbf{r}_{i-1}| \geq t \sup_n |\mathbf{A}_n^T \mathbf{r}_{i-1}|$$

- ✦ an **update** of the residual such that

$$\mathbf{r}_i \in \text{span}(\mathbf{A}_n, n \in \cup_{j=1}^i \Gamma_j)$$

- ✦ Assume =

$$\text{ERC}(I) \quad \sup_{n \notin I} \|\mathbf{A}_I^+ \mathbf{A}_n\|_1 < t$$

- ✦ Conclude: given input $\mathbf{b} = \mathbf{A}_I x_I$, the algorithm is guaranteed to only select atoms in I : $\forall i, \Gamma_i \subset I$

ERC: proof

- Selection rule implies

$$\inf_{l \in \Gamma_i} |\mathbf{A}_l^T \mathbf{r}_{i-1}| \geq t \sup_n |\mathbf{A}_n^T \mathbf{r}_{i-1}| = t \|\mathbf{A}^T \mathbf{r}_{i-1}\|_\infty$$

- It is sufficient to prove by induction that

$$\sup_{n \notin I} |\mathbf{A}_n^T \mathbf{r}_{i-1}| = \|\mathbf{A}_{I^c}^T \mathbf{r}_{i-1}\|_\infty < t \|\mathbf{A}^T \mathbf{r}_{i-1}\|_\infty$$

since this implies $n \notin I \Rightarrow n \notin \Gamma_i$

- Equivalently, we will just prove by induction

$$\frac{\|\mathbf{A}_{I^c}^T \mathbf{r}_{i-1}\|_\infty}{\|\mathbf{A}_I^T \mathbf{r}_{i-1}\|_\infty} < t$$

ERC: proof (ctd)

- **Lemma** $\sup_{\mathbf{r} \in \text{span}(\mathbf{A}_I)} \frac{\|\mathbf{A}_{I^c}^T \mathbf{r}\|_\infty}{\|\mathbf{A}_I^T \mathbf{r}\|_\infty} = \sup_{n \notin I} \|\mathbf{A}_I^+ \mathbf{A}_n\|_1$

- **Proof:**

$$\begin{aligned}
 & \mathbf{r} = \mathbf{A}_{I^c} \mathbf{c} & \mathbf{c} &= (\mathbf{A}_I^T \mathbf{A}_I)^{-1} \mathbf{d} \\
 & \downarrow & & \downarrow \\
 \frac{\|\mathbf{A}_{I^c}^T \mathbf{r}\|_\infty}{\|\mathbf{A}_I^T \mathbf{r}\|_\infty} &= \frac{\|\mathbf{A}_{I^c}^T \mathbf{A}_{I^c} \mathbf{c}\|_\infty}{\|\mathbf{A}_I^T \mathbf{A}_{I^c} \mathbf{c}\|_\infty} &= \frac{\|\mathbf{A}_{I^c}^T \mathbf{A}_I (\mathbf{A}_I^T \mathbf{A}_I)^{-1} \mathbf{d}\|_\infty}{\|\mathbf{d}\|_\infty} \\
 & \leq \|\mathbf{A}_{I^c}^T \mathbf{A}_I (\mathbf{A}_I^T \mathbf{A}_I)^{-1}\|_{\infty \rightarrow \infty} \\
 & = \|(\mathbf{A}_I^T \mathbf{A}_I)^{-1} \mathbf{A}_I^T \mathbf{A}_{I^c}\|_{1 \rightarrow 1} = \|\mathbf{A}_I^+ \mathbf{A}_{I^c}\|_{1 \rightarrow 1}
 \end{aligned}$$

ERC: proof (end)

- For $i=1$ $\mathbf{r}_{i-1} \in \text{span}(\mathbf{A}_I)$
- By induction, if this holds true at step i we obtain

$$\frac{\|\mathbf{A}_I^T \mathbf{r}_{i-1}\|_\infty}{\|\mathbf{A}_I^T \mathbf{r}_{i-1}\|_\infty} \leq \sup_{n \notin I} \|\mathbf{A}_I^+ \mathbf{A}_n\|_1 < t$$

- It follows that $\Gamma_i \subset I$
- By the update rule

$$\mathbf{r}_i \in \text{span}(\mathbf{A}_n, n \in \cup_{j=1}^i \Gamma_j) \subset \text{span}(\mathbf{A}_I)$$

Exact recovery: L_p minimization

Null space

- Null space = kernel

$$z \in \mathcal{N}(\mathbf{A}) \Leftrightarrow \mathbf{A}z = \mathbf{0}$$

- Particular solution vs general solution
 - ♦ particular solution

$$\mathbf{A}x = \mathbf{b}$$

- ♦ general solution

$$\mathbf{A}x' = \mathbf{b} \Leftrightarrow x' - x \in \mathcal{N}(\mathbf{A})$$

Exact recovery: necessary condition

- Notations
 - ◆ index set I
 - ◆ vector z
 - ◆ restriction $z_I = (z_i)_{i \in I}$
- Assume there exists $z \in \mathcal{N}(\mathbf{A})$ with
$$\|z_I\|_f > \|z_{I^c}\|_f$$
- Define $\mathbf{b} := Az_I = A(-z_{I^c})$
- The vector z_I is supported in I but is *not* the minimum norm representation of \mathbf{b}

Exact recovery: sufficient condition

- Assume quasi-triangle inequality

$$\forall x, y \|x + y\|_f \leq \|x\|_f + \|y\|_f$$

- Consider x with support set I and x' with $\mathbf{A}x' = \mathbf{A}x$
- Denote $z := x' - x \in \mathcal{N}(\mathbf{A})$ and observe

$$\begin{aligned} \|x'\|_f &= \|x + z\|_f = \|(x + z)_I\|_f + \|(x + z)_{I^c}\|_f \\ &= \|x + z_I\|_f + \|z_{I^c}\|_f \\ &\geq \|x\|_f - \|z_I\|_f + \|z_{I^c}\|_f \end{aligned}$$

- Conclude:

If $\|z_{I^c}\|_f > \|z_I\|_f$ when $z \in \mathcal{N}(\mathbf{A})$ then I is recoverable

Recoverable supports : the “Null Space Property” (I)

- **Theorem I** [*Donoho & Huo 2001 for l_1 , G. & Nielsen 2003 for L_p*]
 - ♦ Assumption 1: sub-additivity (for quasi-triangle inequality)

$$f(a + b) \leq f(a) + f(b), \forall a, b$$

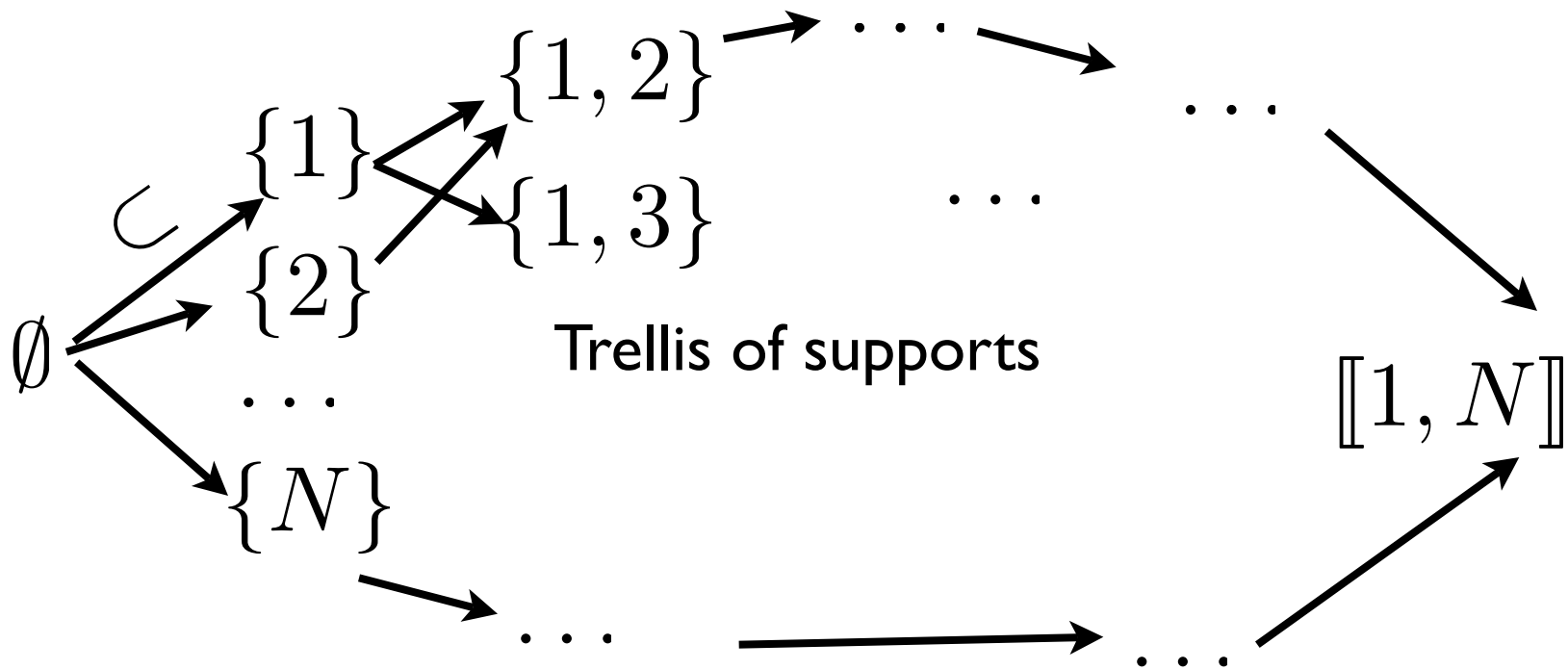
- ♦ Assumption 2:

NSP

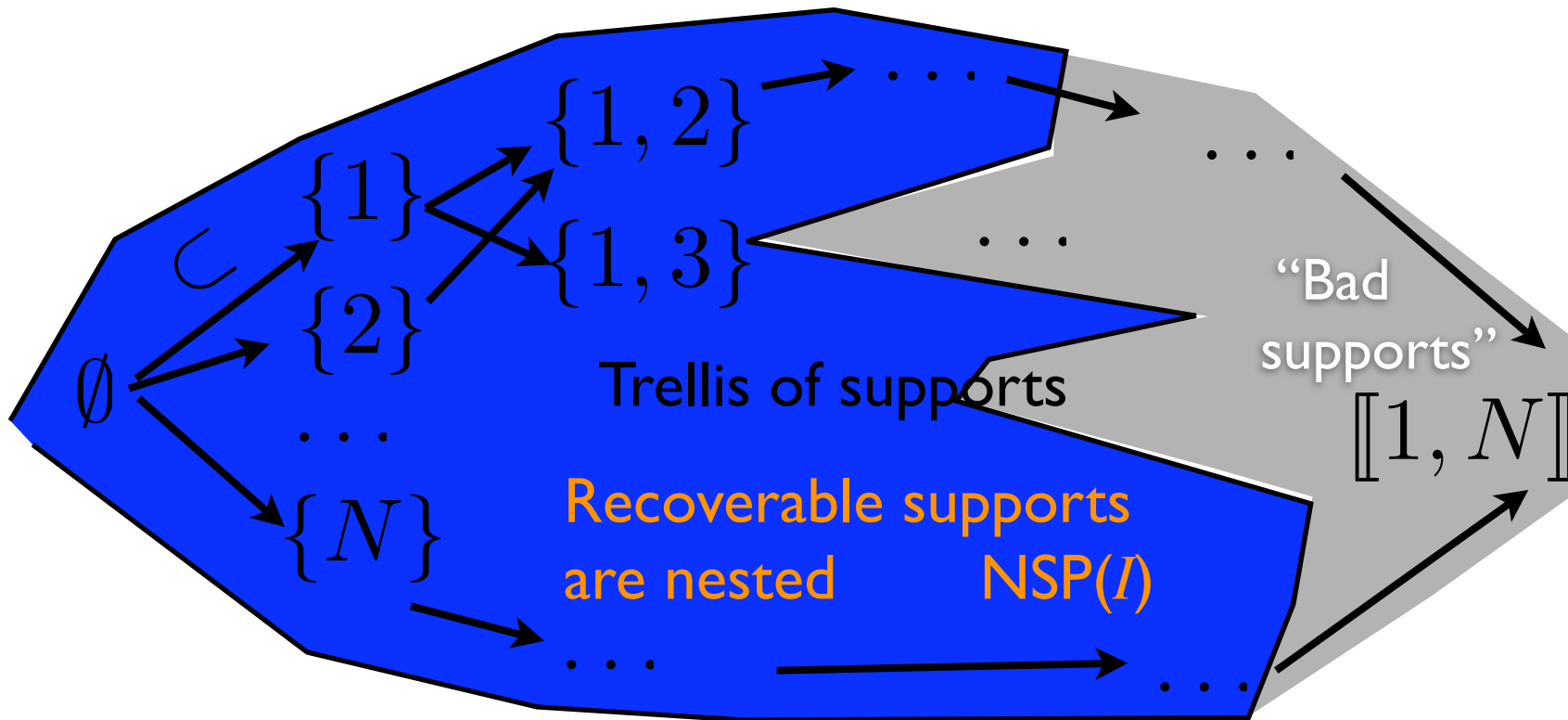
$$\|z_I\|_f < \|z_{I^c}\|_f \text{ when } z \in \mathcal{N}(\mathbf{A}), z \neq 0$$

- ♦ Conclusion: x_f^* recovers every x supported in I
- ♦ The result is sharp: if NSP fails on support I there is **at least one failing vector** x supported in I

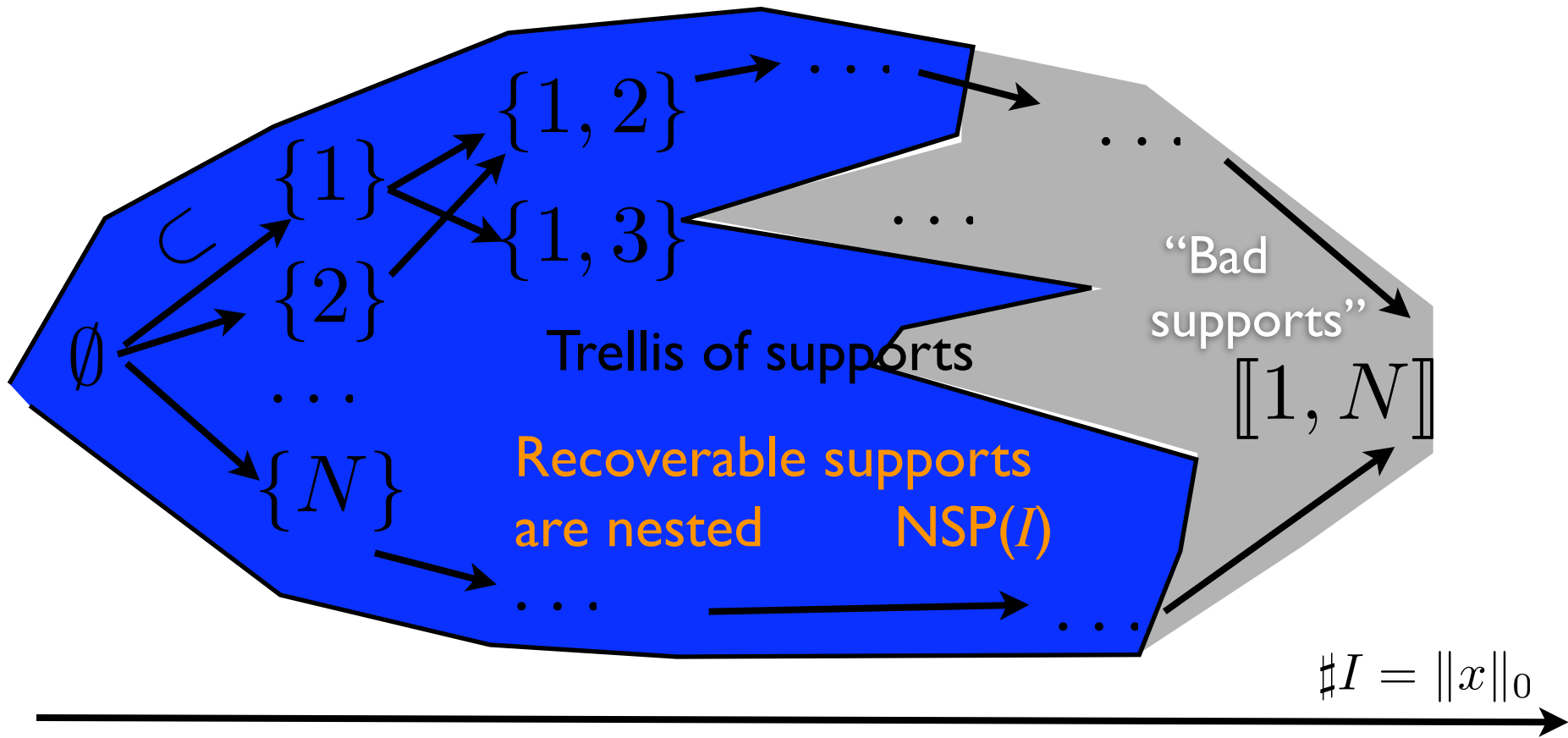
From “recoverable” supports to “sparse” vectors



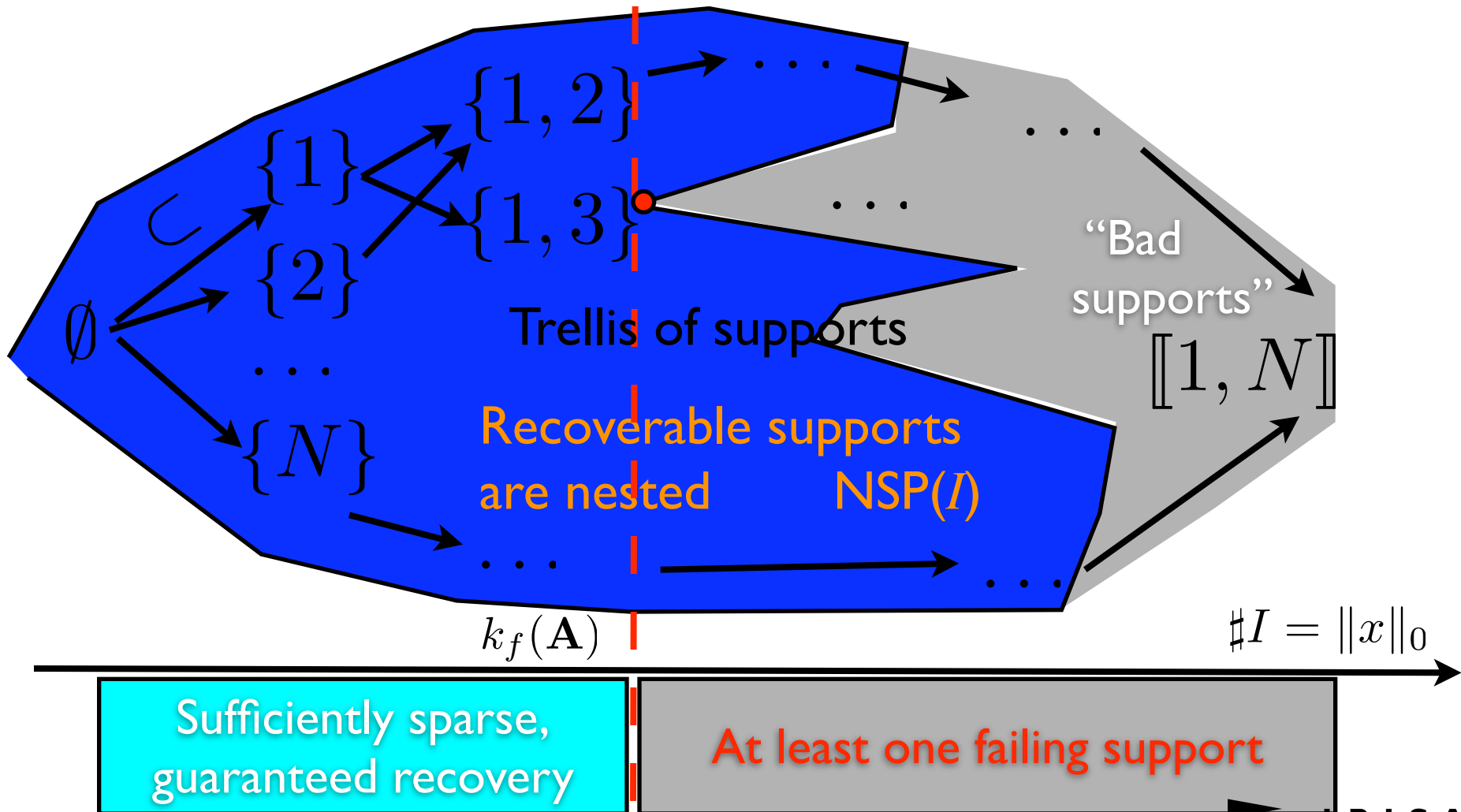
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Recoverable sparsity levels: the “Null Space Property” (2)

- Corollary 1 [*Donoho & Huo 2001 for l_1 , G. Nielsen 2003 for l_p*]

- ♦ Definition :

$I_k =$ index of k largest components of z

- ♦ Assumption :

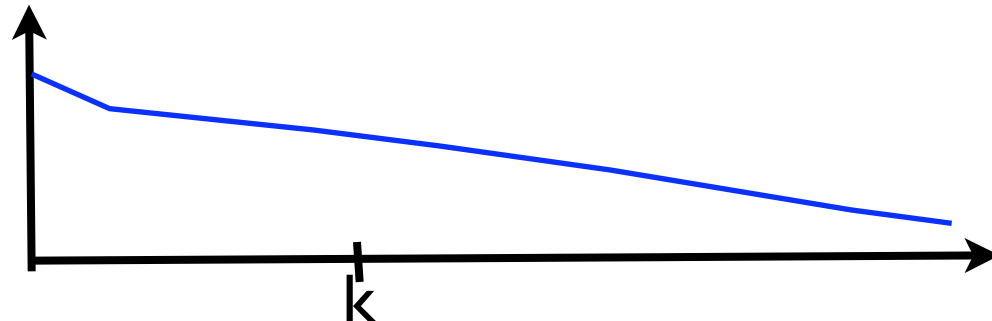
NSP

$$\|z_{I_k}\|_f < \|z_{I_k^c}\|_f \quad \text{when } z \in \mathcal{N}(\mathbf{A}), z \neq 0$$

- ♦ Conclusion: x_f^* recovers every x with $\|x\|_0 \leq k$
- ♦ The result is sharp: if NSP fails there is **at least one failing vector** x with $\|x\|_0 = k$

Interpretation of NSP

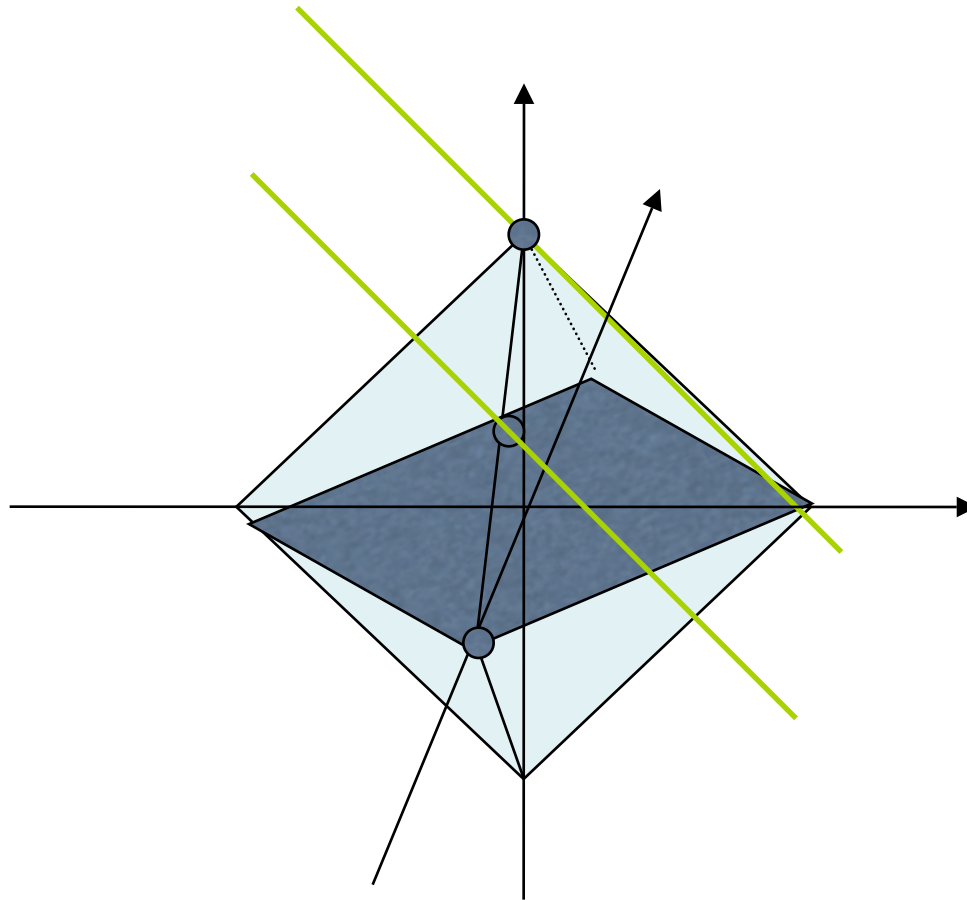
- Geometry in coefficient space:
 - ◆ consider an element z of the Null Space of A
 - ◆ order its entries in decreasing order



- ◆ the mass of the largest k -terms should not exceed that of the tail $\|z_{I_k}\|_f < \|z_{I_k^c}\|_f$

All elements of the null space must be rather “flat”

Geometric picture



Greedy vs LI

ERC implies NSP

$$\text{ERC}(I) \quad \sup_{n \notin I} \|\mathbf{A}_I^+ \mathbf{A}_n\|_1 < t$$

If columns of \mathbf{A}_I
linearly independent



$$\text{NSP}(I, \ell^1, t) \quad \|z_I\|_1 \leq t \cdot \|z_{I^c}\|_1 \quad \text{when } z \in \mathcal{N}(\mathbf{A}), z \neq 0$$

ERC implies NSP: proof

- Assume ERC(I) + linear independence $\sup_{n \notin I} \|\mathbf{A}_I^+ \mathbf{A}_n\|_1 < t$
- Consider $z \in \mathcal{N}(\mathbf{A})$
 - ♦ write $\mathbf{A}_I z_I = -\mathbf{A}_{I^c} z_{I^c}$
 - ♦ apply pseudo-inverse $z_I = -\mathbf{A}_I^+ \mathbf{A}_{I^c} z_{I^c}$

$$\|z_I\|_1 \leq \|\mathbf{A}_I^+ \mathbf{A}_{I^c}\|_{1 \rightarrow 1} \cdot \|z_{I^c}\|_1 \leq t \|z_{I^c}\|_1$$

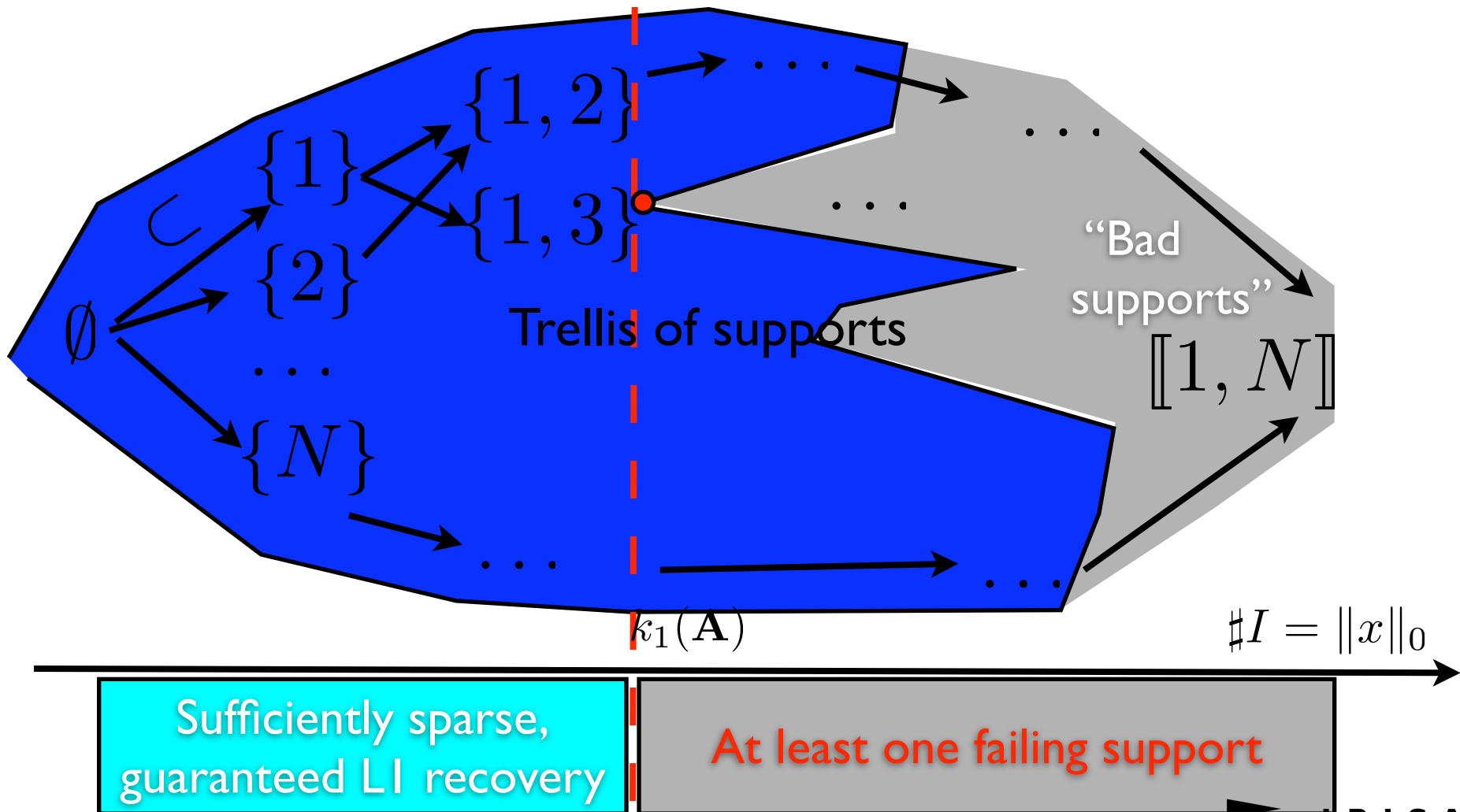
Greedy vs LI: summary

- If *MP is guaranteed to recover all vectors with support I , where the atoms in I are linearly independent, then LI has the same guarantee
- If *MP recovers all k -sparse vectors, then LI has the same guarantee

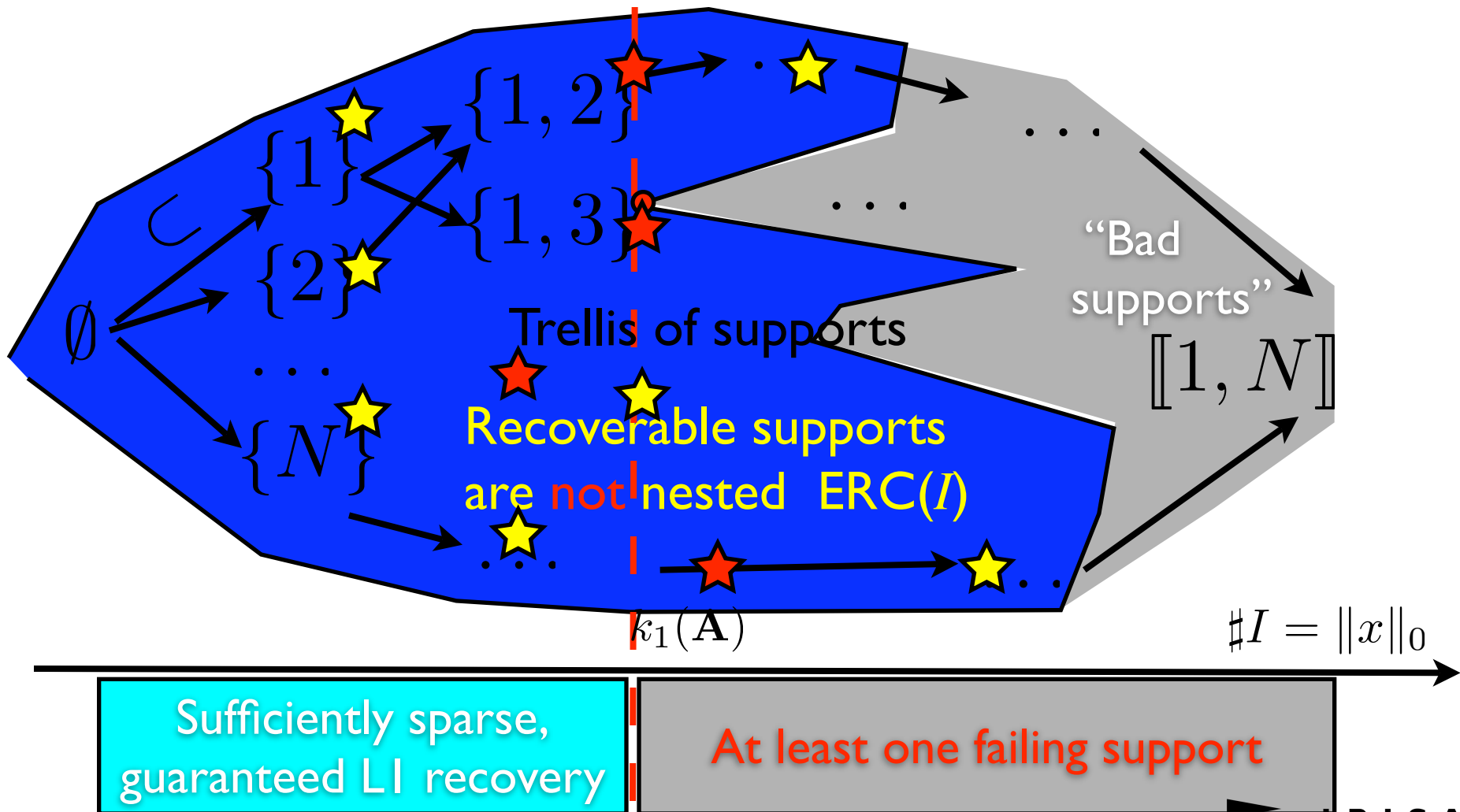
$$k_{*MP}(\mathbf{A}) \leq k_1(\mathbf{A}), \forall \mathbf{A}$$

- Warning: there are support sets I
 - ◆ not recovered by LI, while recovered by MP
 - ◆ ... but columns of \mathbf{A}_I are linearly dependent.
 - ◆ Example: $I = \llbracket 1, N \rrbracket$ when \mathbf{A} is $m \times N$

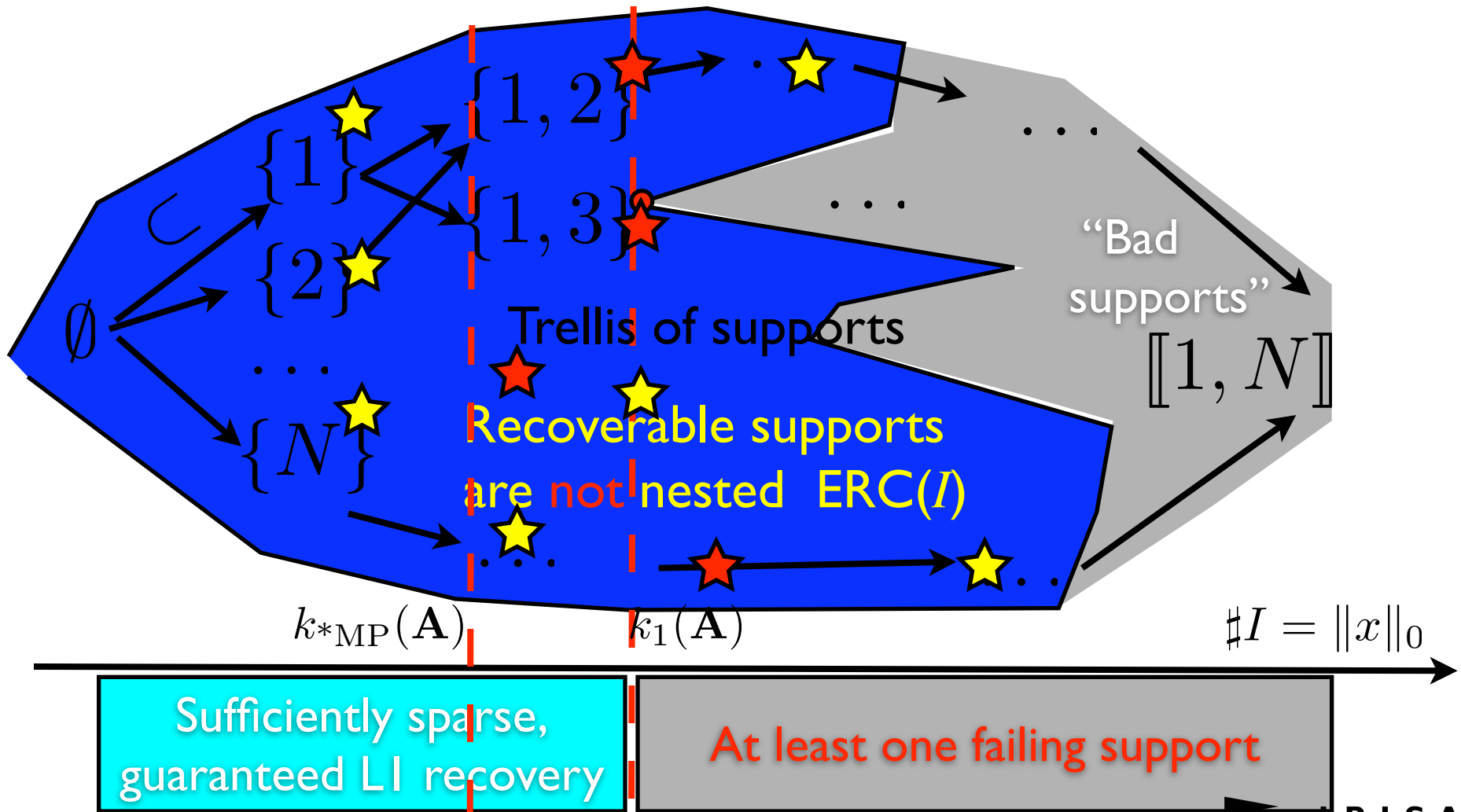
Greedy vs LI: summary



Greedy vs LI: summary



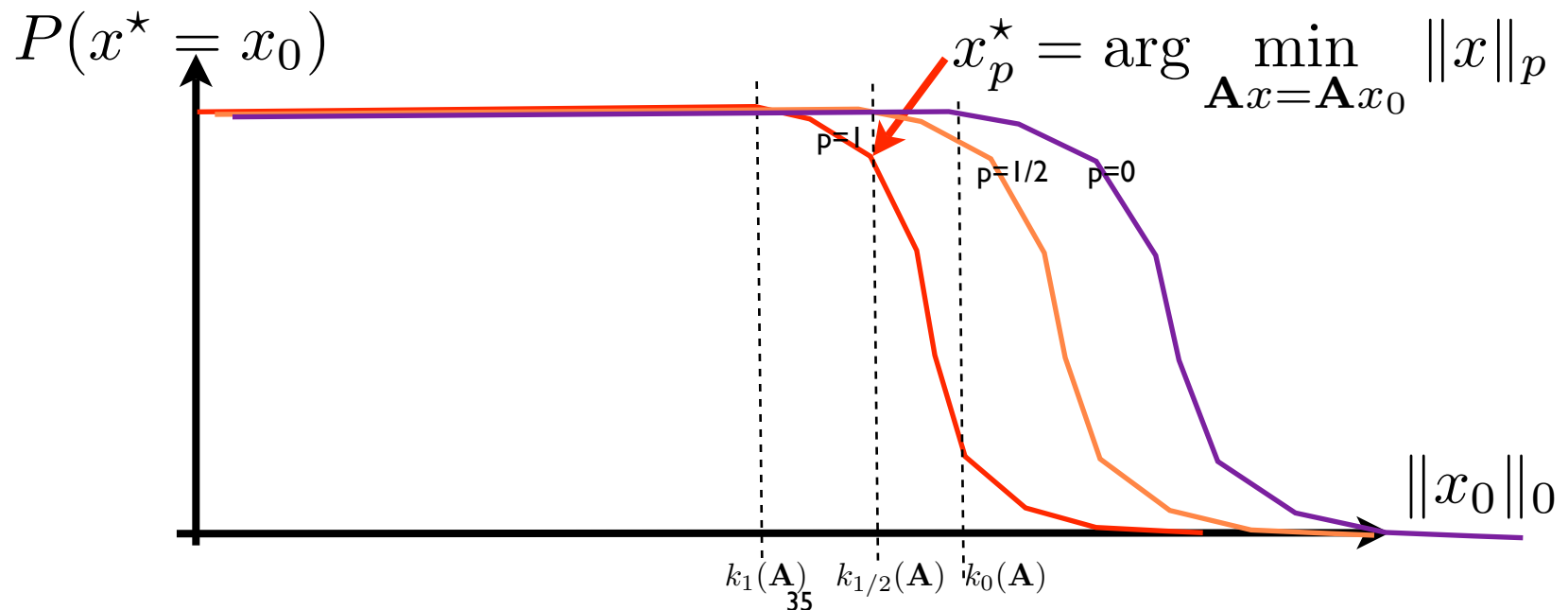
Greedy vs LI: summary



LI vs Lp

Critical sparsity levels for different L_p norms

Back to empirical observations + extrapolation :



Do we always have $k_1(\mathbf{A}) \leq k_p(\mathbf{A}) \leq k_0(\mathbf{A}), 0 \leq p \leq 1$?

Lp better than L1 (I)

- **Theorem 2** [G. Nielsen 2003]

- ◆ Assumption 1: **sub-additivity** of sparsity measures f, g

$$f(a + b) \leq f(a) + f(b), \forall a, b$$

- ◆ Assumption 2: the function $t \mapsto \frac{f(t)}{g(t)}$ is **non-increasing** ↘

- ◆ Conclusion: $k_g(\mathbf{A}) \leq k_f(\mathbf{A}), \forall \mathbf{A}$

Minimizing $\|x\|_f$ can recover vectors which are less sparse than required for guaranteed success when minimizing $\|x\|_g$

Lp better than L1 (2)

● Example

◆ sparsity measures $f(t) = t^p, g(t) = t^q, 0 \leq p \leq q \leq 1$

◆ sub-additivity

$$|a + b|^p \leq |a|^p + |b|^p, \forall a, b, 0 \leq p \leq 1$$

◆ function $\frac{f(t)}{g(t)} = t^{p-q}$ is non-increasing

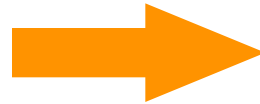
◆ therefore

$$k_1(\mathbf{A}) \leq k_q(\mathbf{A}) \leq k_p(\mathbf{A}) \leq k_0(\mathbf{A}), \forall \mathbf{A}$$

Lp better than L1: proof

- 1) Since f/g non-decreasing:

$$z_1 \geq z_2 \geq 0$$



$$\frac{f(z_1)}{g(z_1)} \leq \frac{f(z_2)}{g(z_2)}$$

- 2) Similarly

$$z_1 \geq \dots \geq z_N \geq 0$$



$$\frac{\|z_{1:k}\|_f}{\|z_{1:k}\|_g} \leq \frac{\|z_{k+1:N}\|_f}{\|z_{k+1:N}\|_g}$$

$$I_k = 1 : k$$

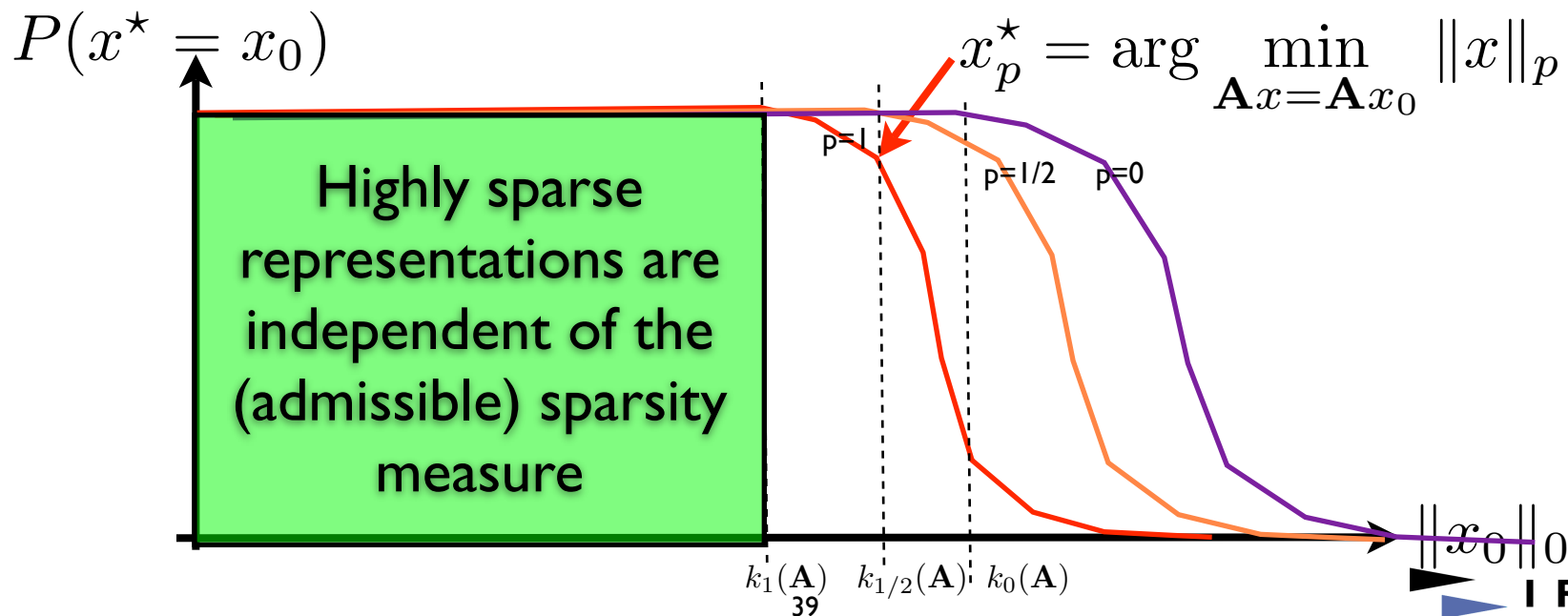
$$I_k^c = k + 1 : N$$

$$\frac{\|z_{I_k}\|_f}{\|z_{I_k^c}\|_f} \leq \frac{\|z_{I_k}\|_g}{\|z_{I_k^c}\|_g}$$

- 3) Conclusion : if $\text{NSP}(g,t,k)$ then $\text{NSP}(f,t,k)$

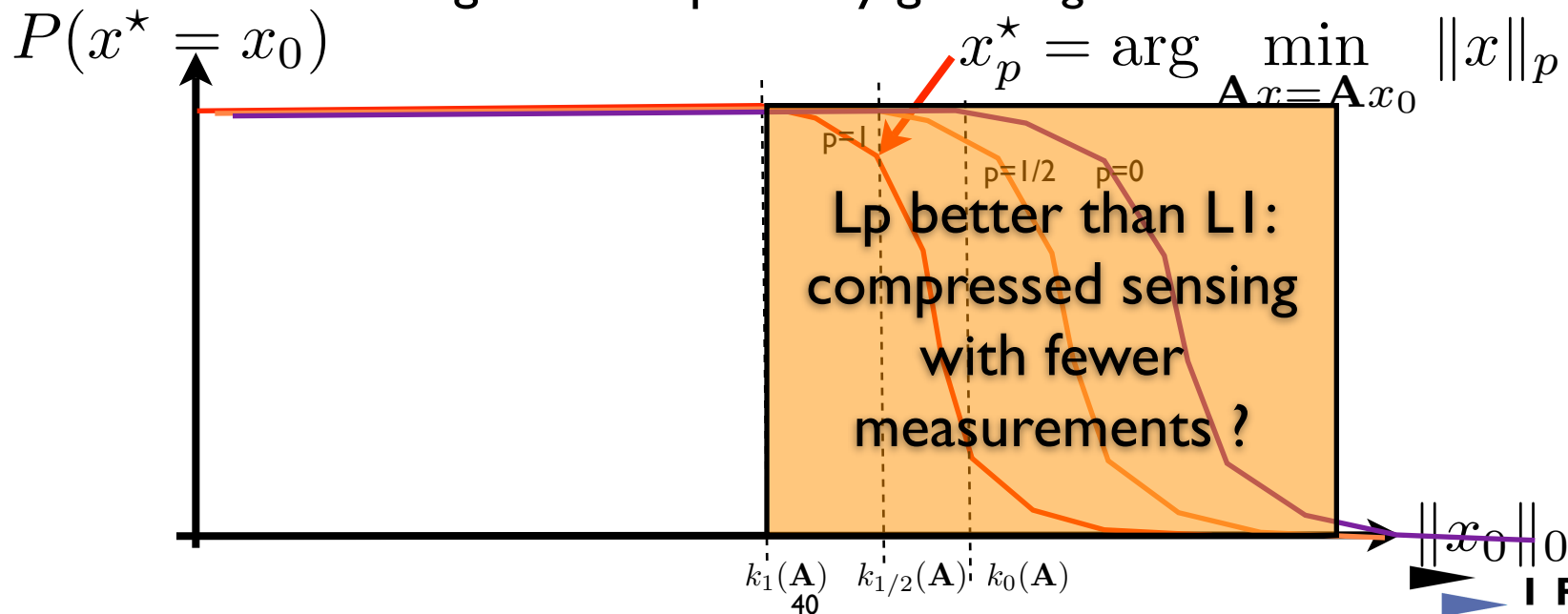
Lp better than L1 (2)

- At sparsity levels where L1 is guaranteed to “succeed”, all Lp $p \leq 1$ is also guaranteed to succeed



Lp better than L1 (3)

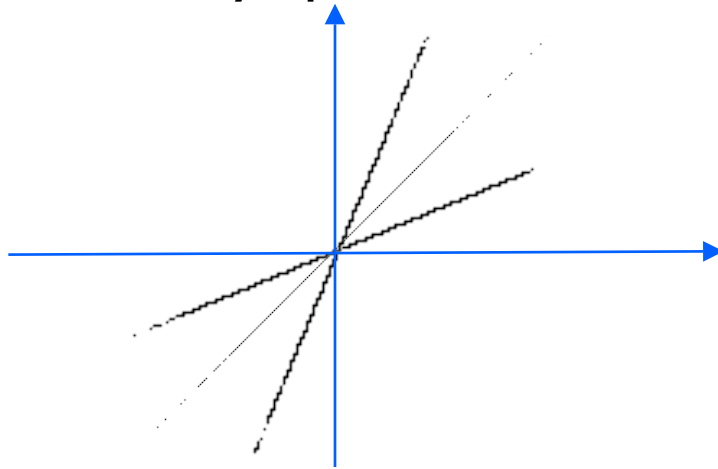
- + Lp $p < 1$ can succeed where L1 fails
 - ◆ How much improvement? Quantify $k_p(\mathbf{A})$?
- - Lp $p < 1$: nonconvex, has many local minima
 - ◆ Better recovery with Lp *principle*
 - ◆ Challenge : actual provably good *algorithms*?



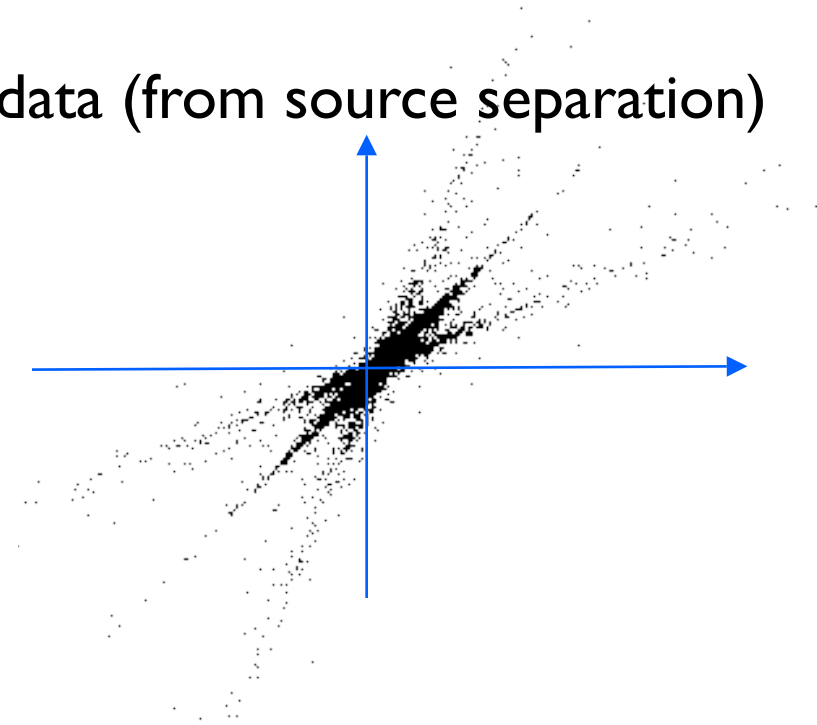
Stability and robustness

Stability

Exactly sparse data



Real data (from source separation)



Stability

- Exact recovery: $\mathbf{b} = \mathbf{A}x$
 - ◆ sparsity assumption $\|x\|_0 \leq k_p(\mathbf{A}) < m$
 - ◆ recovery: $x_p^*(\mathbf{b}) = x$

- Stability: relax sparsity assumption
 - ◆ best k-term approximation

$$\sigma_k(x) = \inf_{\|y\|_0 \leq k} \|x - y\|$$

- ◆ goal = stable recovery = instance optimality

$$\|x_p^*(\mathbf{b}) - x\| \leq C \cdot \sigma_k(x)$$

Instance optimality for L_p minimization

- Assumption:

$$\overset{\text{NSP}(k, \ell^p, t)}{\|z_{I_k}\|_p^p \leq t \cdot \|z_{I_k^c}\|_p^p} \quad \text{when } z \in \mathcal{N}(\mathbf{A}), z \neq 0$$

- Conclusion: instance optimality for all x

$$\|x_p^*(\mathbf{b}) - x\|_p^p \leq C(t) \cdot \sigma_k(x)_p^p$$

$$C(t) := 2 \frac{1+t}{1-t}$$

Robustness

- Noiseless model $\mathbf{b} = \mathbf{A}x$
- Noisy model $\mathbf{b} = \mathbf{A}x + \mathbf{e}$
 - ◆ measurement noise
 - ◆ modeling error
 - ◆ numerical inaccuracies ...
- Goal: robust estimation

$$\|x_p^*(\mathbf{b}) - x\| \leq C\|e\| + C'\sigma_k(x)$$

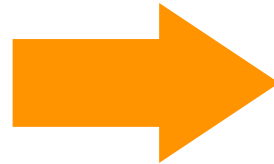
- Tool: restricted isometry property

Restricted Isometry Property

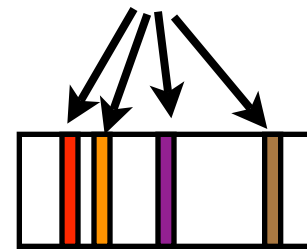
- Definition



N columns



$$n \in I, \#I \leq k$$



max over



A_I subsets I

$$\frac{N!}{k!(N-k)!}$$

- Computation ?

- ✦ naively: combinatorial
- ✦ **open question:** NP ? NP-complete ?

$$\delta_k := \sup_{\#I \leq k, c \in \mathbb{R}^k} \left| \frac{\|A_I c\|_2^2}{\|c\|_2^2} - 1 \right|$$

RIP, stability, robustness

RIP(k, δ)

$$\delta_{2k}(\mathbf{A}) \leq \delta$$

[Candès 2008]



$$t := \sqrt{2}\delta / (1 - \delta)$$

NSP(k, ℓ^1, t)

$$\|z_{I_k}\|_1 \leq t \cdot \|z_{I_k^c}\|_1 \quad \text{when } z \in \mathcal{N}(\mathbf{A}), z \neq 0$$

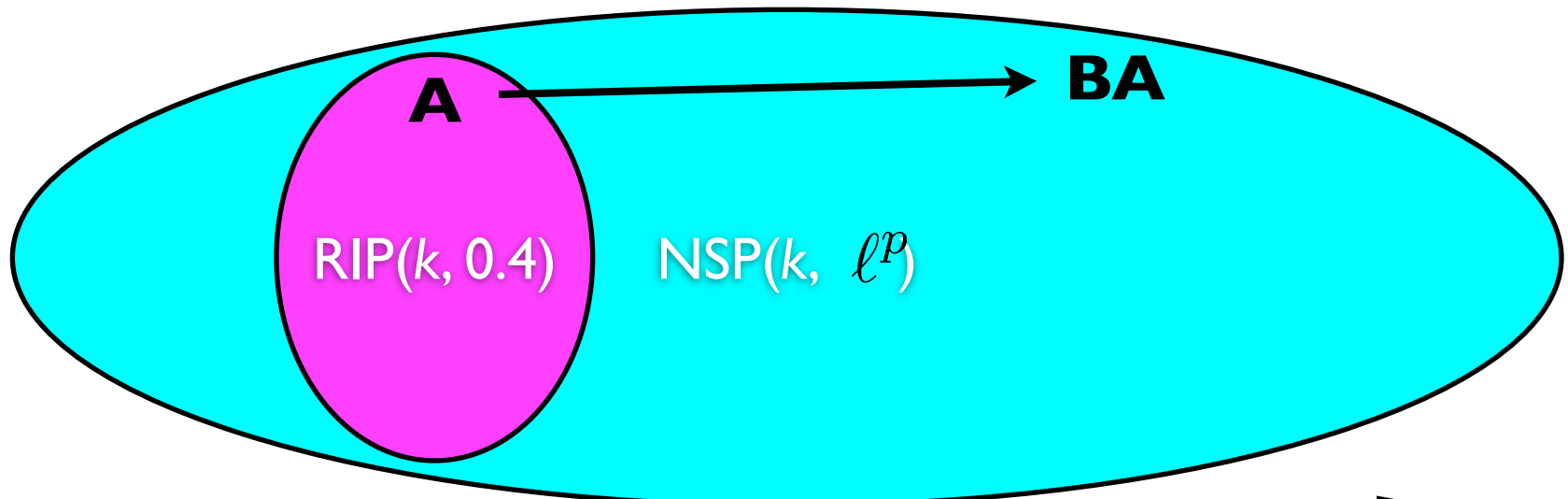
- LI-recovery of k -sparse vectors, **robust to noise**

when $\mathbf{b} = \mathbf{A}x + \mathbf{n}$ $\delta_{2k}(\mathbf{A}) < \sqrt{2} - 1 \approx 0.414$

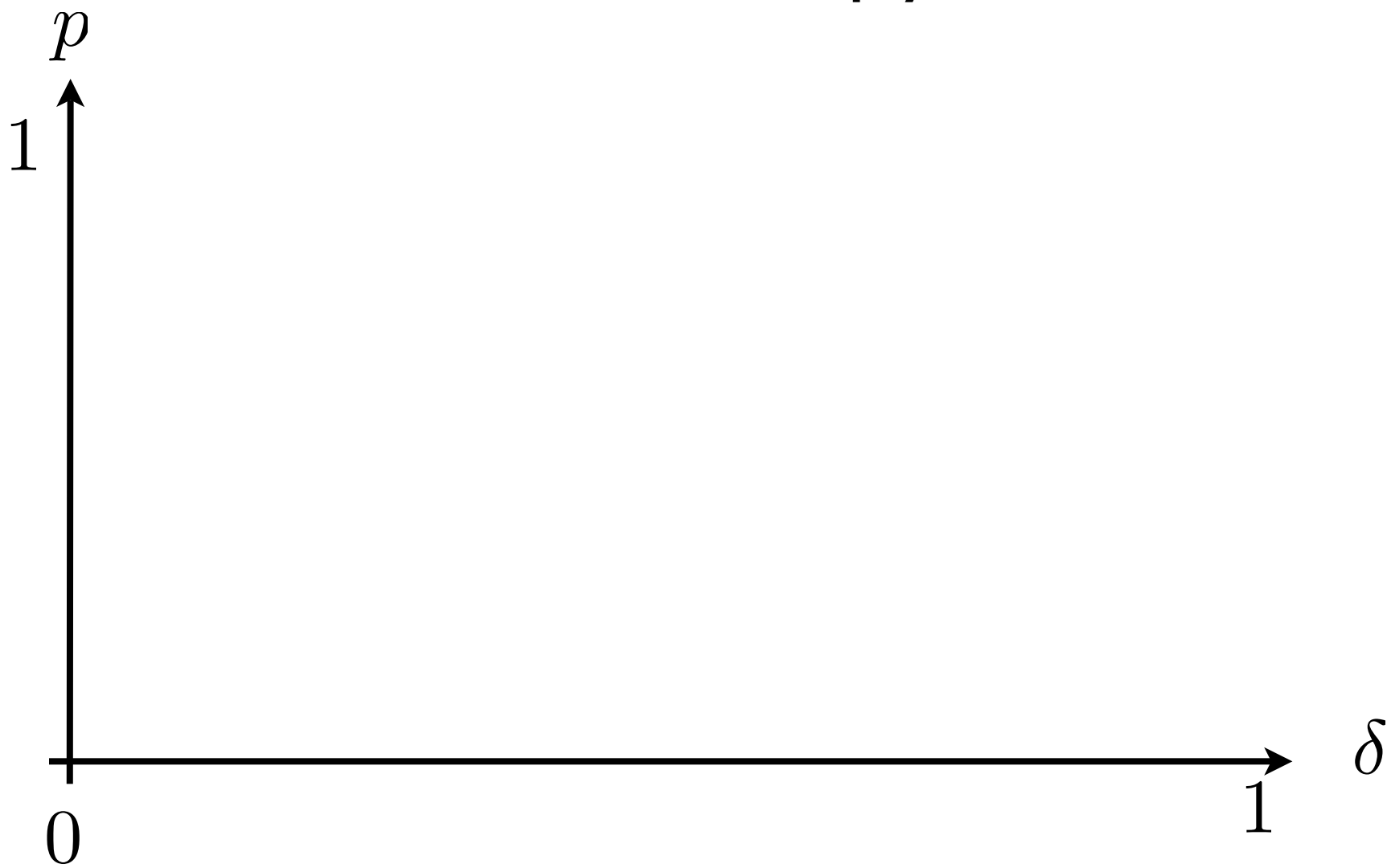
- ✦ Foucart-Lai 2008: ℓ_p with $p < 1$, and $\delta_{2k}(\mathbf{A}) < 0.4531$
- ✦ Chartrand 2007, Saab & Yilmaz 2008: other RIP condition for $p < 1$
- ✦ G., Figueras & Vandergheynst 2006: robustness with f -norms

How sharp is the RIP condition ?

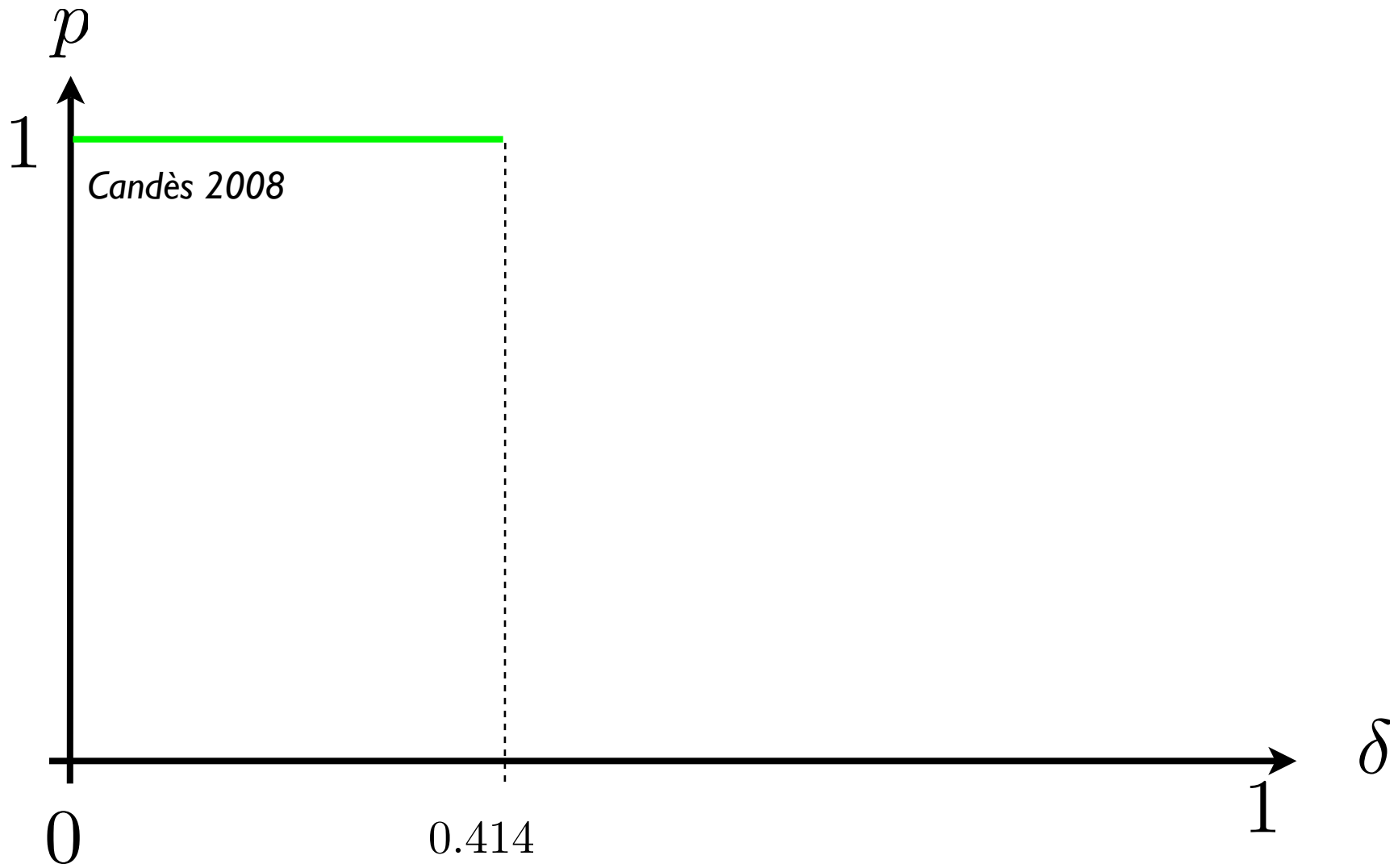
- The Null Space Property for L_p
 - ♦ “algebraic” + sharp property, only depends on $\mathcal{N}(\mathbf{A})$
 - ♦ invariant by linear transforms $\mathbf{A} \rightarrow \mathbf{B}\mathbf{A}$
- The $\text{RIP}(k, \delta)$ condition is “metric” ... and not invariant
 - ♦ even with “rescaled” RIP



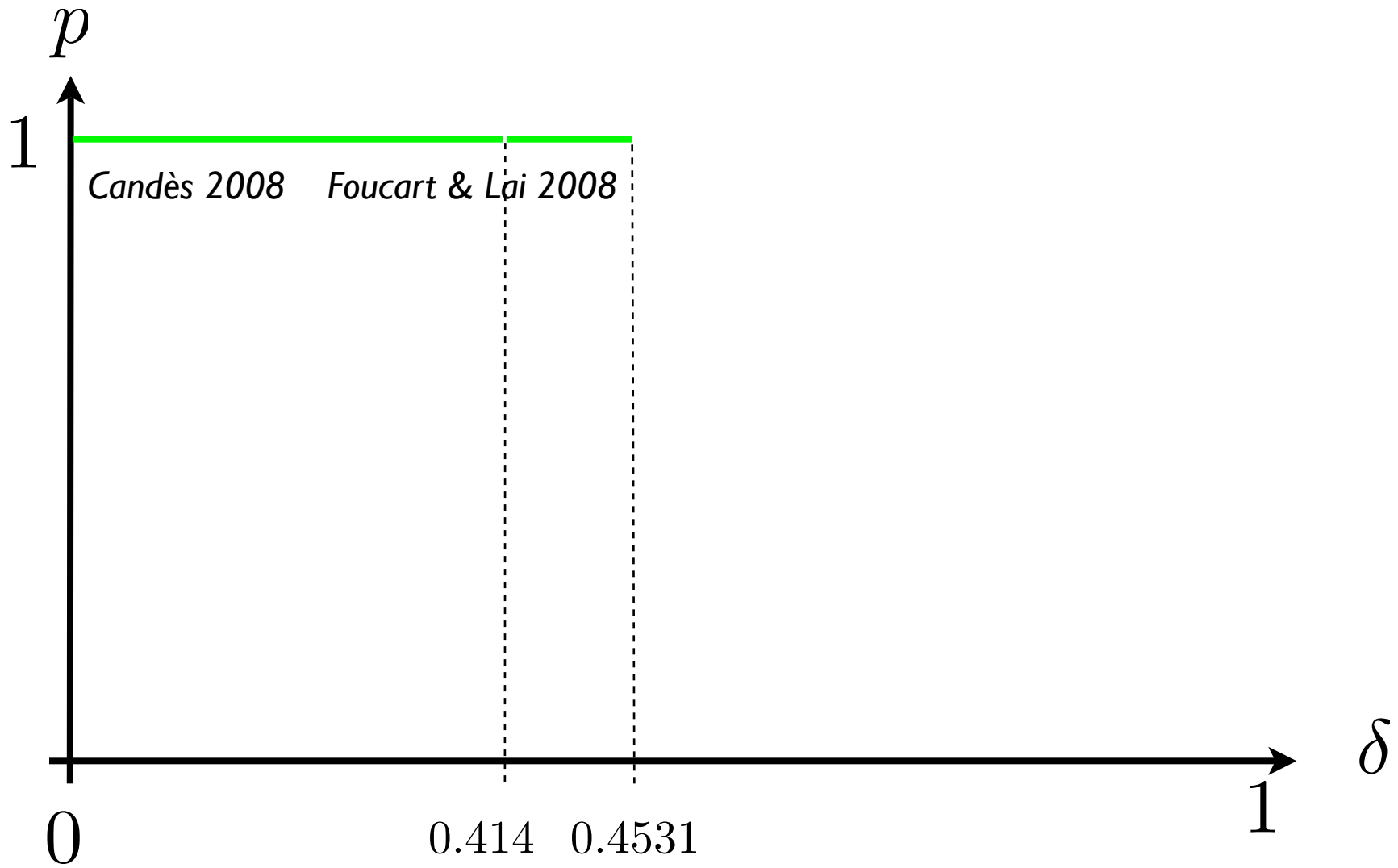
When does $\delta_{2k}(\mathbf{A}) < \delta$ imply $k \leq k_p(\mathbf{A})$?



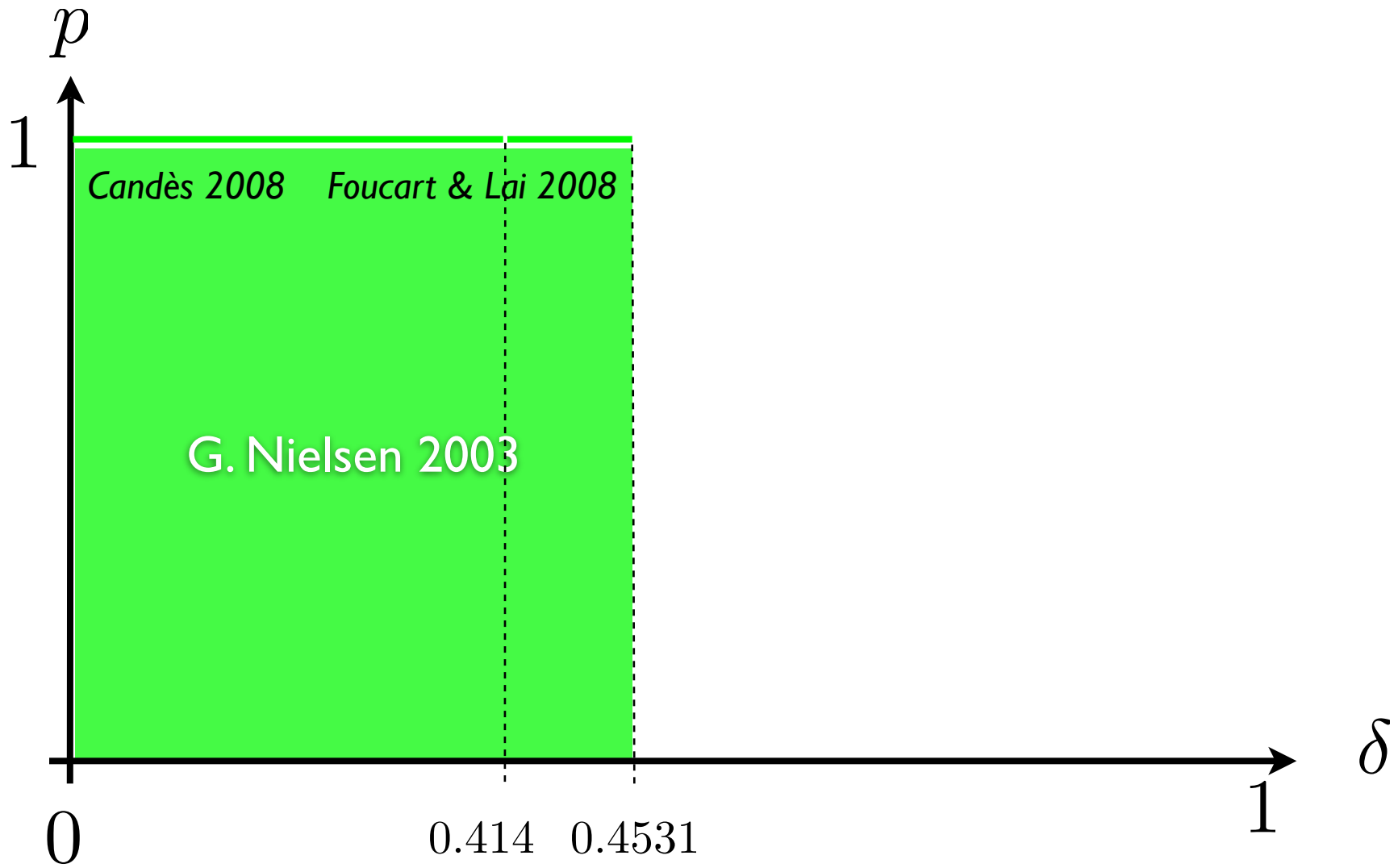
When does $\delta_{2k}(\mathbf{A}) < \delta$ imply $k \leq k_p(\mathbf{A})$?



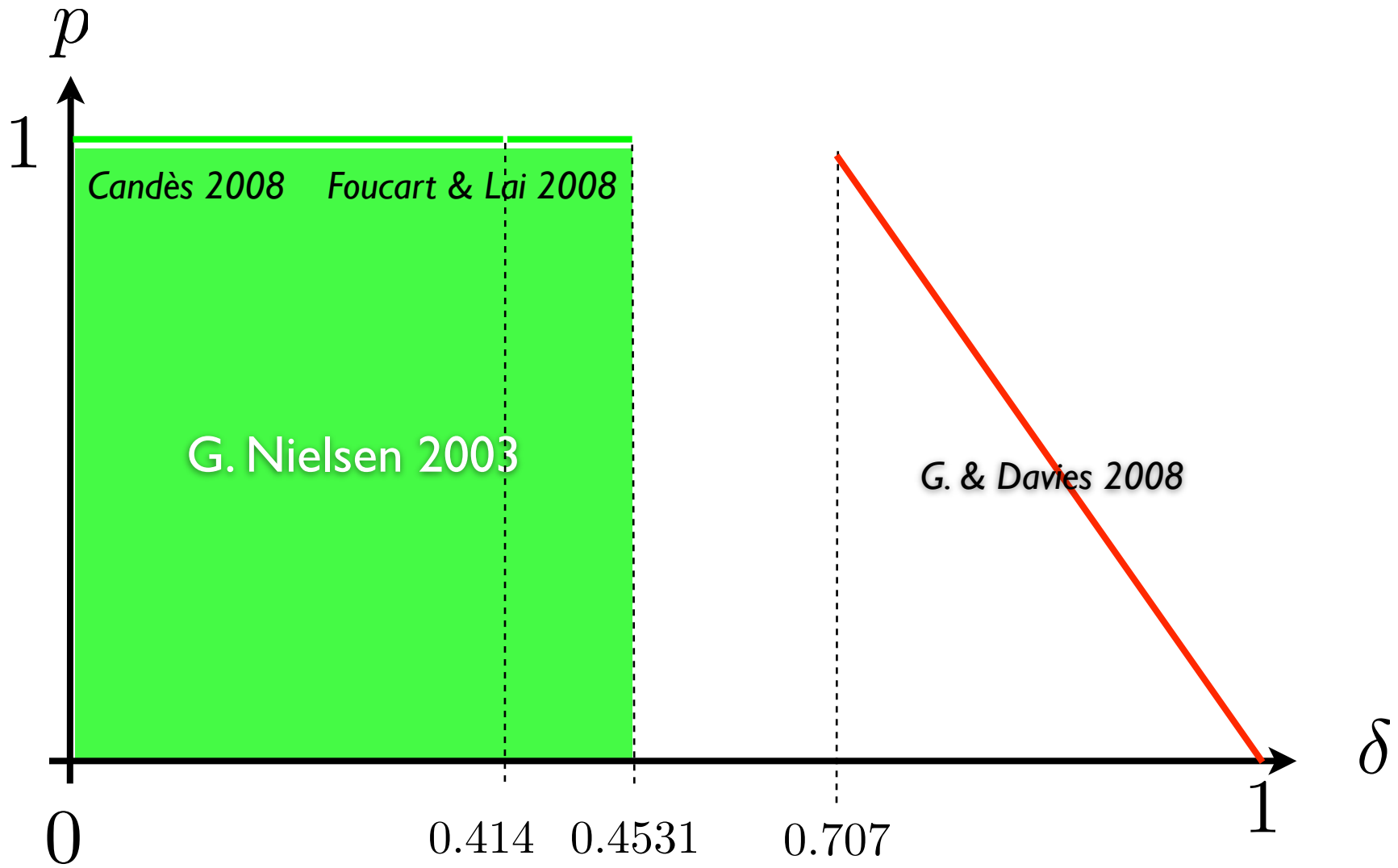
When does $\delta_{2k}(\mathbf{A}) < \delta$ imply $k \leq k_p(\mathbf{A})$?



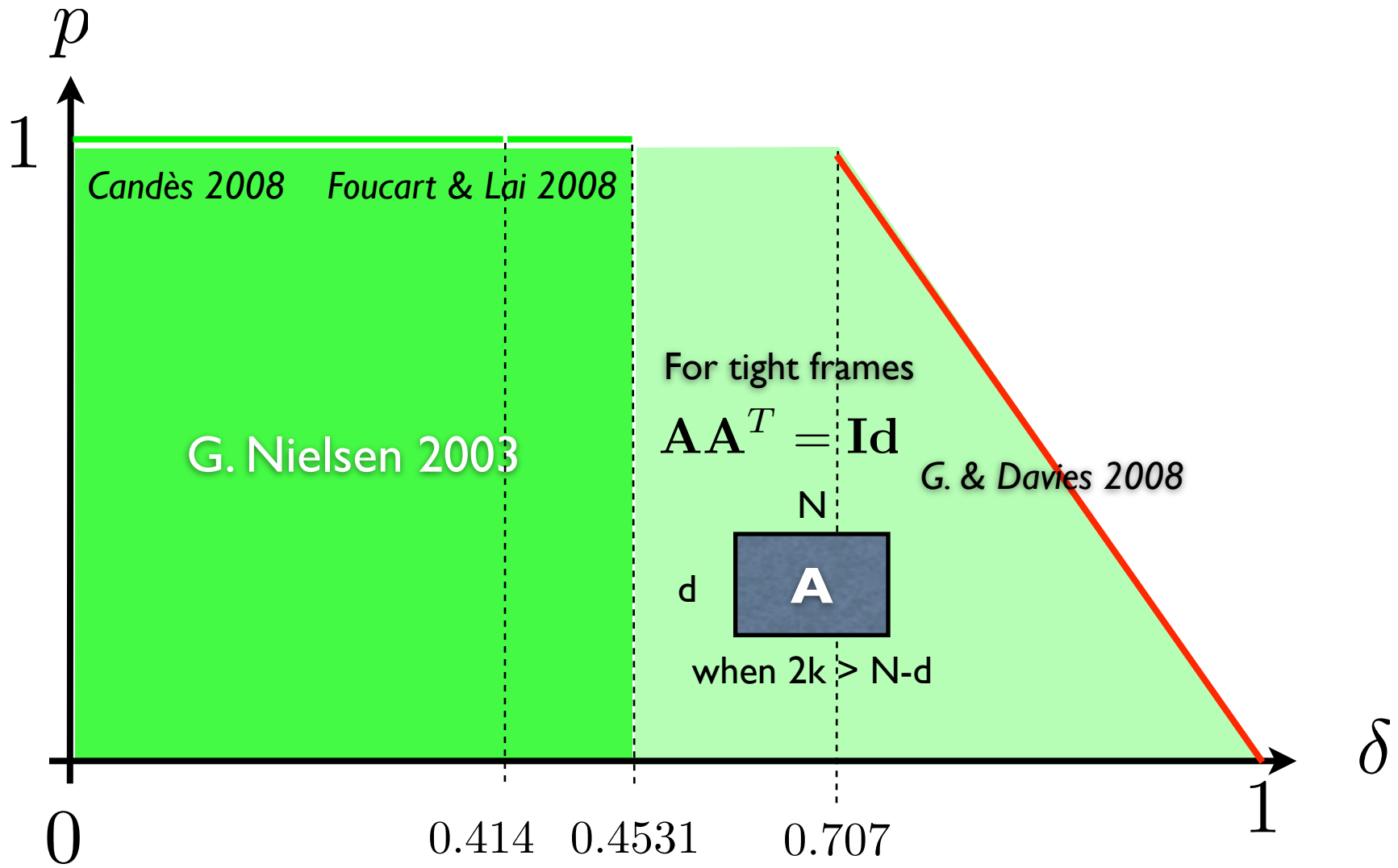
When does $\delta_{2k}(\mathbf{A}) < \delta$ imply $k \leq k_p(\mathbf{A})$?



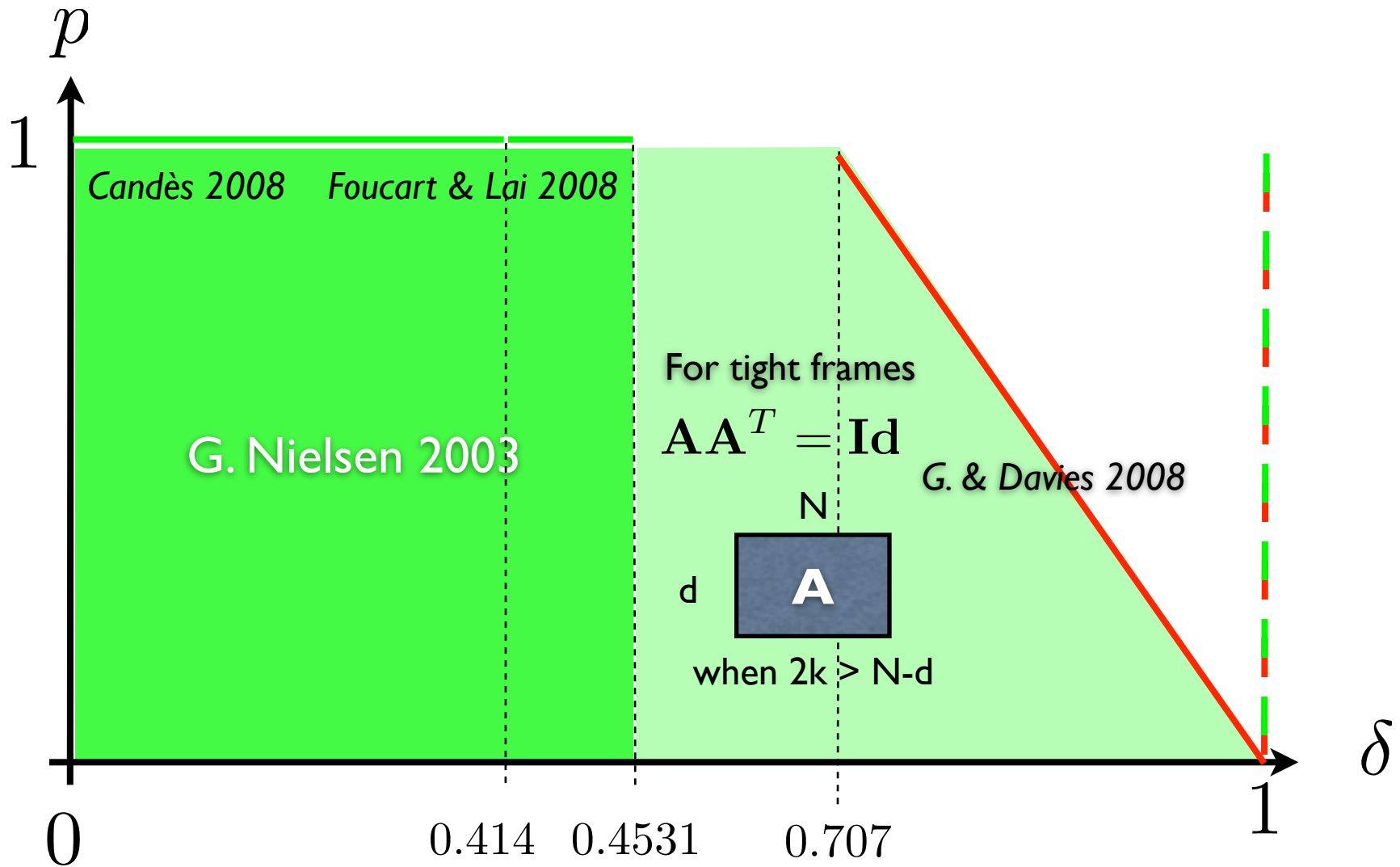
When does $\delta_{2k}(\mathbf{A}) < \delta$ imply $k \leq k_p(\mathbf{A})$?



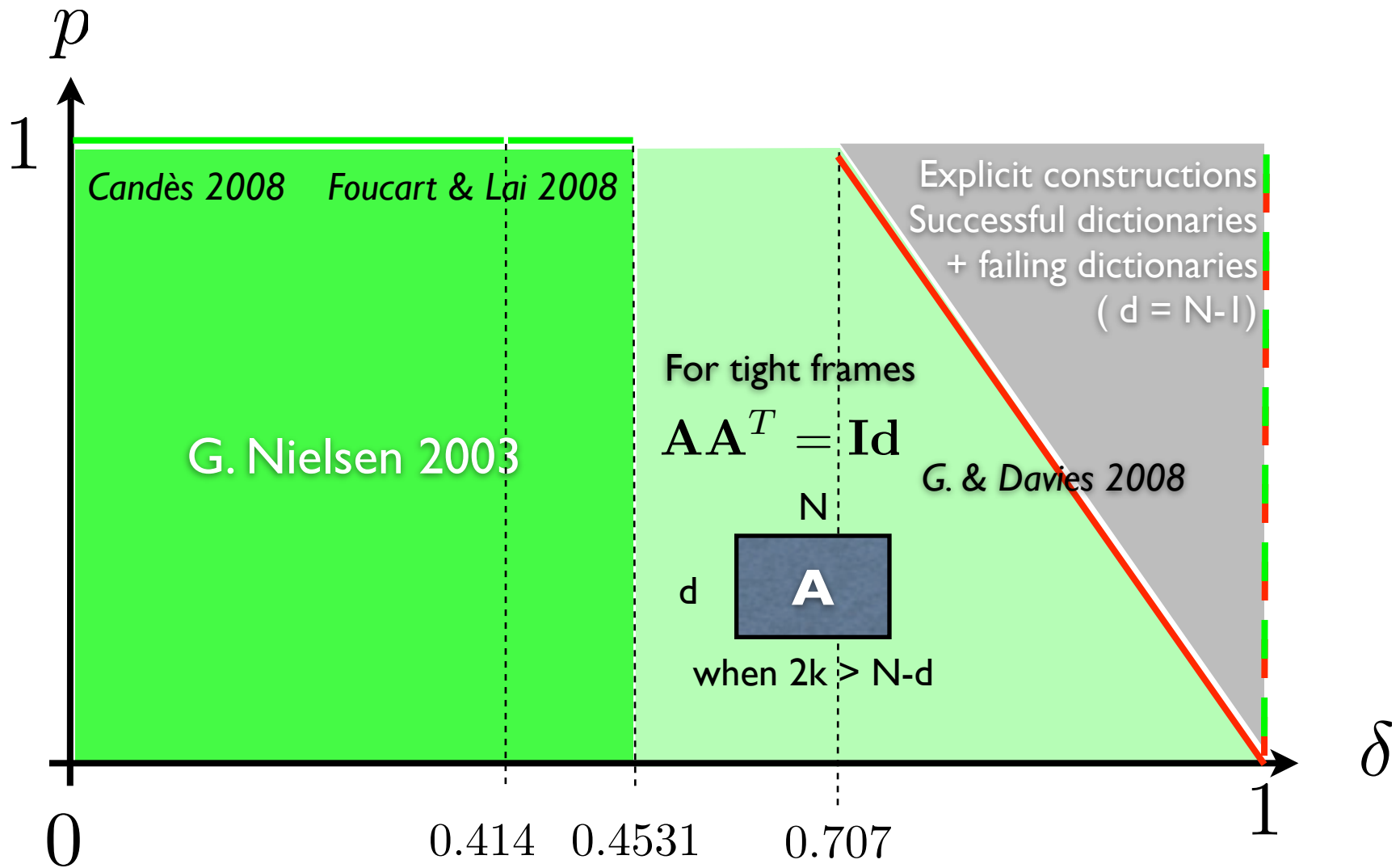
When does $\delta_{2k}(\mathbf{A}) < \delta$ imply $k \leq k_p(\mathbf{A})$?



When does $\delta_{2k}(\mathbf{A}) < \delta$ imply $k \leq k_p(\mathbf{A})$?



When does $\delta_{2k}(\mathbf{A}) < \delta$ imply $k \leq k_p(\mathbf{A})$?



Summary

- Recovery conditions based on number of nonzero components $\|x\|_0$

$$k_{\text{MP}}^*(\mathbf{A}) \leq k_1(\mathbf{A}) \leq k_p(\mathbf{A}) \leq k_q(\mathbf{A}) \leq k_0(\mathbf{A}), \forall \mathbf{A}$$

- **Warning:**
 - ◆ there often exists vectors beyond these critical sparsity levels, which are recovered
 - ◆ there often exists vectors beyond these critical sparsity levels, where the successful algorithm is not the one we would expect